Adding Concurrency to Smart Contracts

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Abstract

Modern cryptocurrency systems, such as Ethereum, permit complex financial transactions through scripts called *smart contracts*. These smart contracts are executed many, many times always without real concurrency. First, all smart contracts are executed by *miners*, before appending them to the blockchain. Later, those contracts are serially executed by *validators* to verify that the smart contracts were executed correctly by miners.

Serial execution limits system throughput and fails to exploit today's concurrent multi and cluster architectures. Nevertheless, serial execution appears to be required: contracts and contract programming languages have a serial semantics.

presents a novel way to permit miners and validators to execute smart cont
ues adapted from software transactional memory. Miners execut
llowing non-conflicting contracts to proceed conc
hedule for a block's transactions, This s
m used by validators to re-e
peedup


Bitcoin
Satoshi Nakamoto 2008
Cryptocurrency
No central authority
Anyone can participate

Abstraction: Distributed Ledger

Cash

Date	Description	Increase	Decrease	Balance
Jan. 1, 20X3	Balance forward			$ 50,000
Jan. 2, 20X3	Collected receivable	$ 10,000		60,000
Jan. 3, 20X3	Cash sale	5,000		65,000
Jan. 5, 20X3	Paid rent		7,000	58,000
Jan. 7, 20X3	Paid salary		3,000	55,000
Jan. 7, 20X3	Cash sale	4,000		59,000
Jan. 8, 20X3	Paid bills		2,000	57,000
Jan. 10, 20X3	Paid tax		1,000	56,000
Jan. 12, 20X3	Collected receivable	7,000		63,000

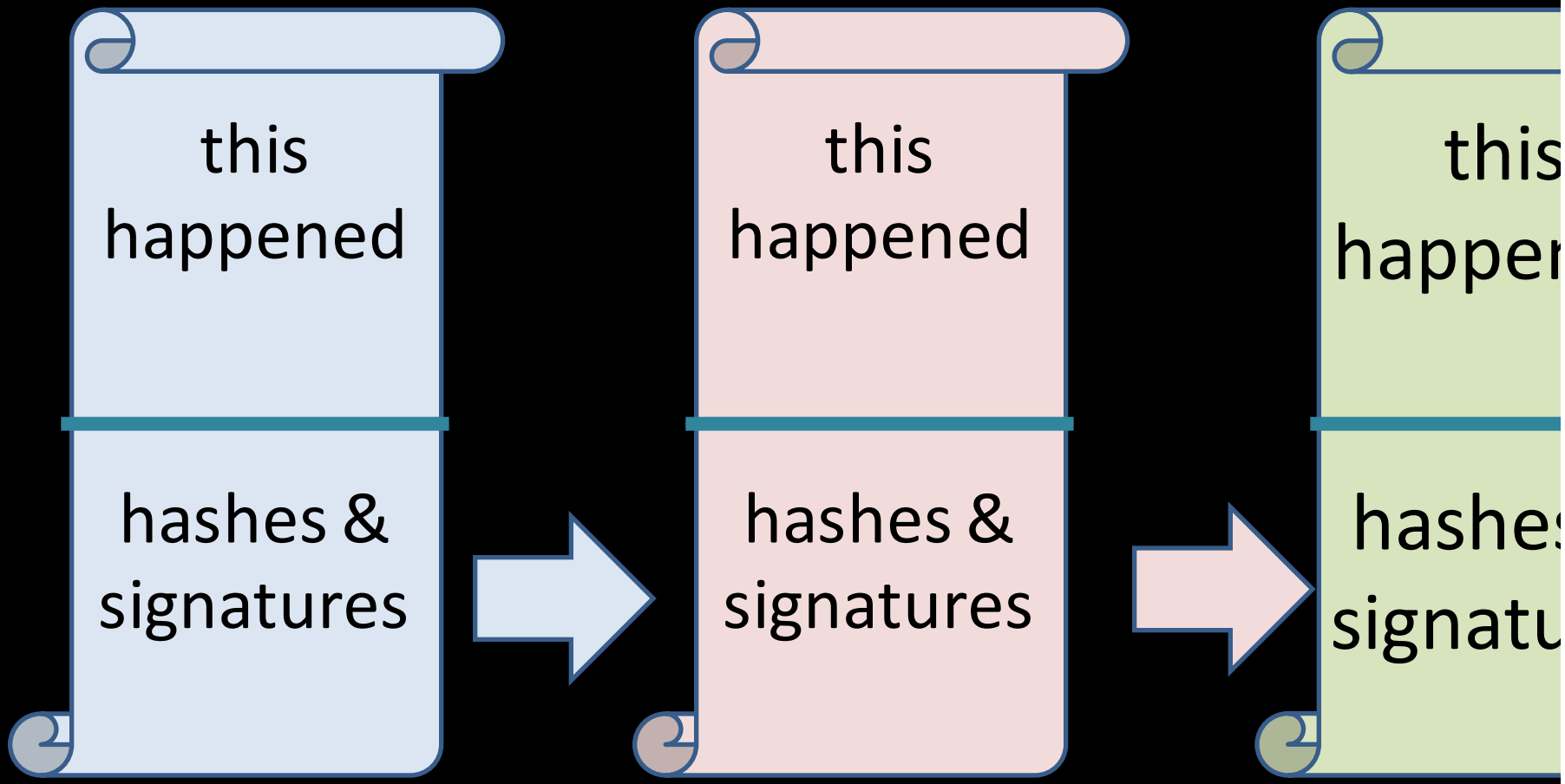
Append-only list of events

Tamper-proof!

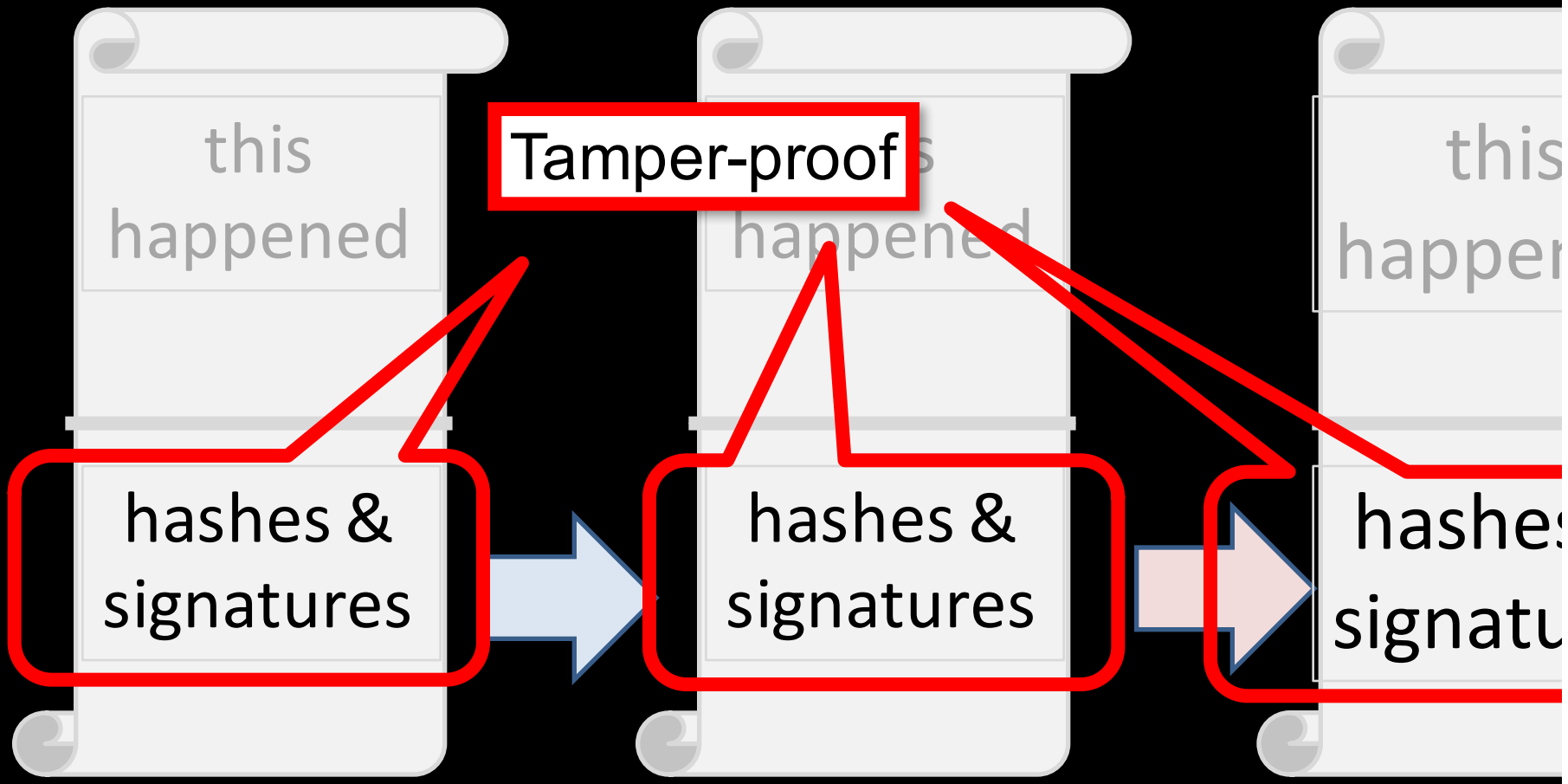
Everyone agrees on content

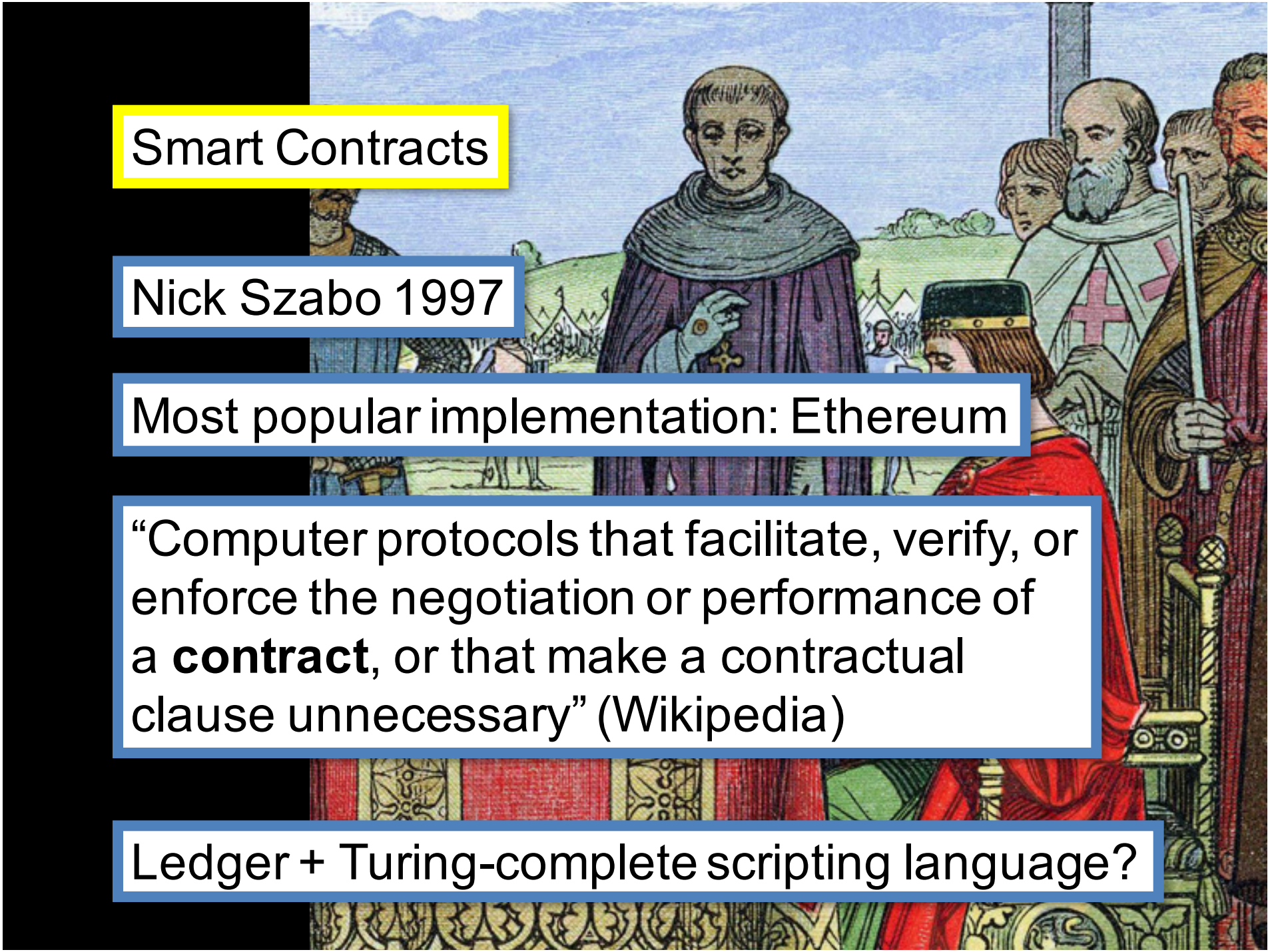
Not just financial

Implementation: Blockchain



Implementation: Blockchain





Smart Contracts

Nick Szabo 1997

Most popular implementation: Ethereum

"Computer protocols that facilitate, verify, or enforce the negotiation or performance of a **contract**, or that make a contractual clause unnecessary" (Wikipedia)

Ledger + Turing-complete scripting language?

```
contract Ballot {
  mapping(address => Voter)
     public voters;
     … // more state decls
  function vote(uint proposal)
    Voter sender = voters[msg.sender];
    if (sender.voted)
      throw;
    sender.voted = true;
    sender.vote = proposal;
    proposals[proposal].voteCount
      += sender.weight;
  }
  …
}
```

Looks like an object in a language

```
contract Ballot {
  mapping(address => Voter)
      public voters;
      … // more state decls
  function vote(uint proposal)
    Voter sender = voters[msg.sender];
    if (sender.voted)
      sender.voted = true;
    }
  …
}
```

Long-lived state

Built-in data types: maps, arrays, scalars.

Tracks who can vote, who voted, choices.

```
contract Ballot {
  mapping(address => Voter)
     public voters;
     … // more state decls
  function vote(uint proposal)
    Voter sender = voters[msg.sender];
    if (sender.voted)
      sender.voted = true;
      proposals[proposal].voteCount
        += sender.weight;
  }
  …
}
```

Functions to manipulate state

Vote for a particular proposal

```
contract Ballot {
  mapping(address => Voter)
     public voters;
     … // more state decls
  function vote(uint proposal)
    Voter sender = voters[msg.sender];
    if (sender.voted)
      throw;
    sender.voted = true;
    sender.vote = proposal;
    proposals[proposal].voteCount
      += sender.weight;
  }
  …
}
```

No voting twice

```
contract Ballot {
  mapping(address => Voter)
     public voters;
     … // more state decls
  function vote(uint p
    Voter sender = voters[msg.sender];
    if (sender.voted)
      throw;
    sender.voted = true;
    sender.vote = proposal;
    proposals[proposal].voteCount
      += sender.weight;
  }
  …
}
```

Record vote

```
contract Ballot {
  mapping(address => Voter)
     public voters;
     … // more state decls
  function vote(uint proposal)
    Voter sender = voters[msg.sender];
    if (sender.voted)
    sender.voted = true;
    sender.vote = proposal;
    proposals[proposal].voteCount
      += sender.weight;
  }
  …
}
```

On a blockchain this is a shared object!

All contract code executed *sequentially*

Every transaction executed sequentially by *everyone*

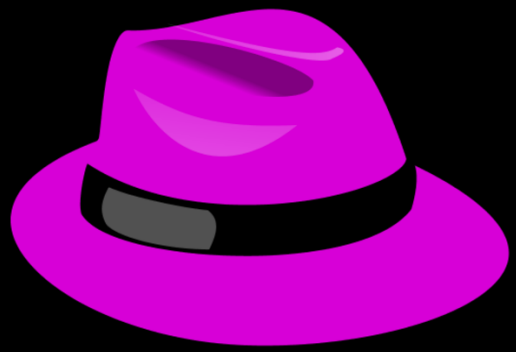
No concurrency control built in to contract language

Big idea #1: permit parallel execution, adapting STM techniques, i.e., speculative execution with rollback

Big idea #2: publish concurrent schedules to the blockchain for everyone to exploit parallelism

Smart Contracts on the Blockchain

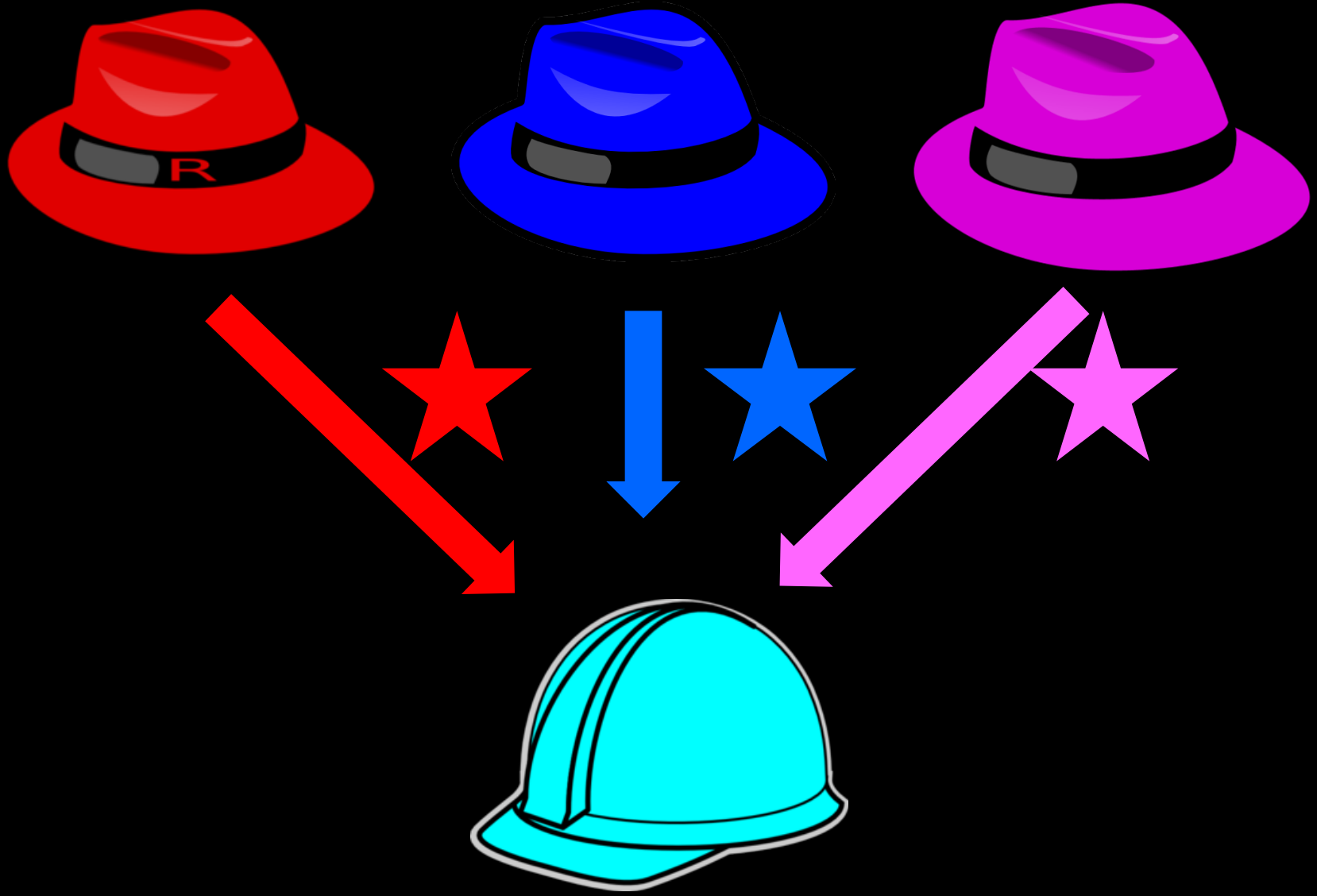
Clients ….



Miners …

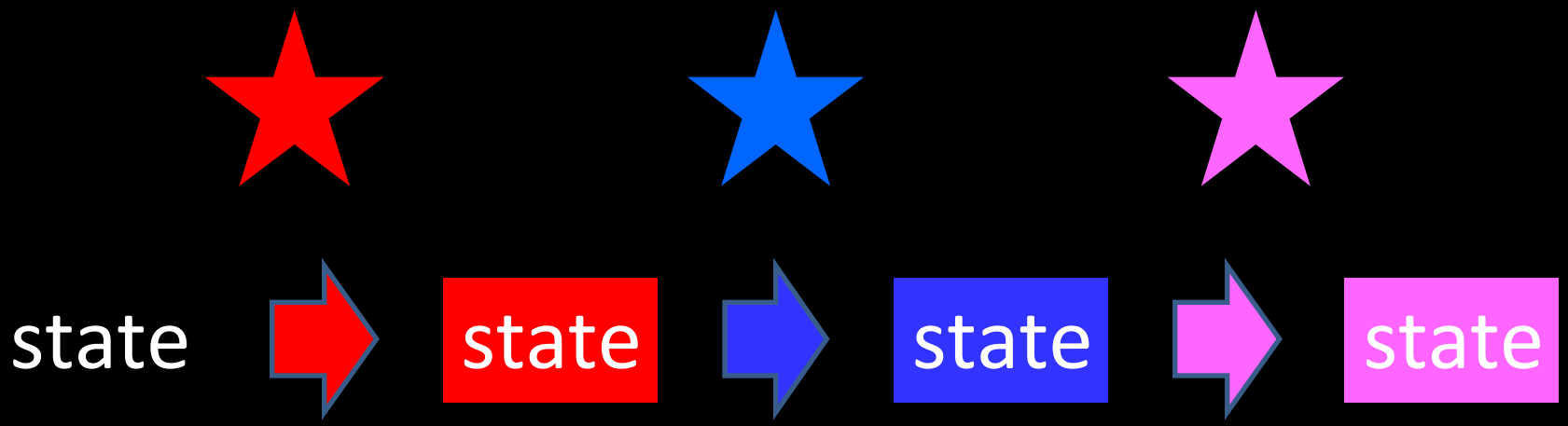
Validators …

Clients send transactions & contracts to miners

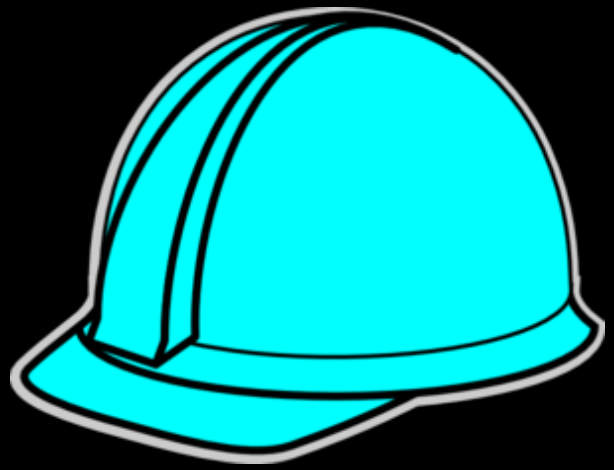


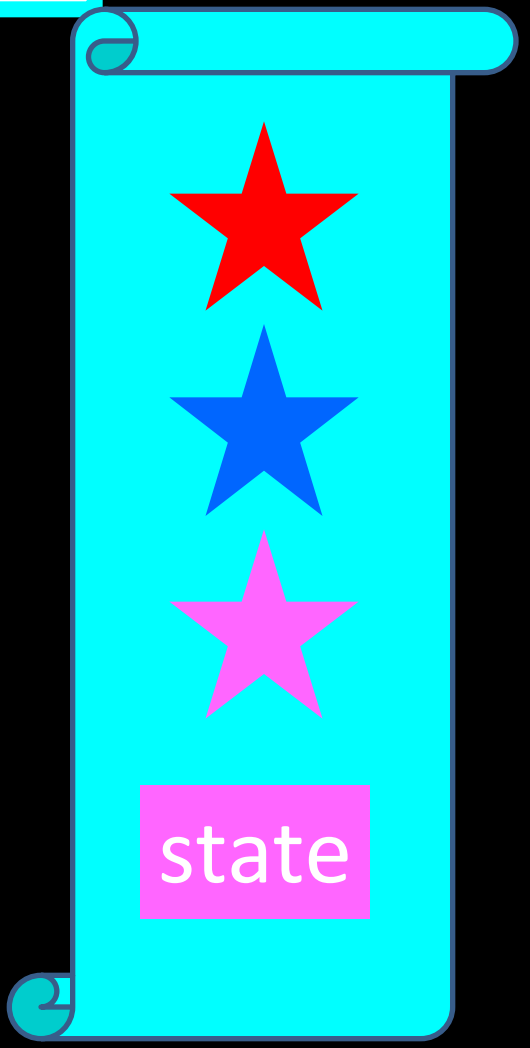
Miners collect transactions…

Apply them one-at-a-time to compute new state

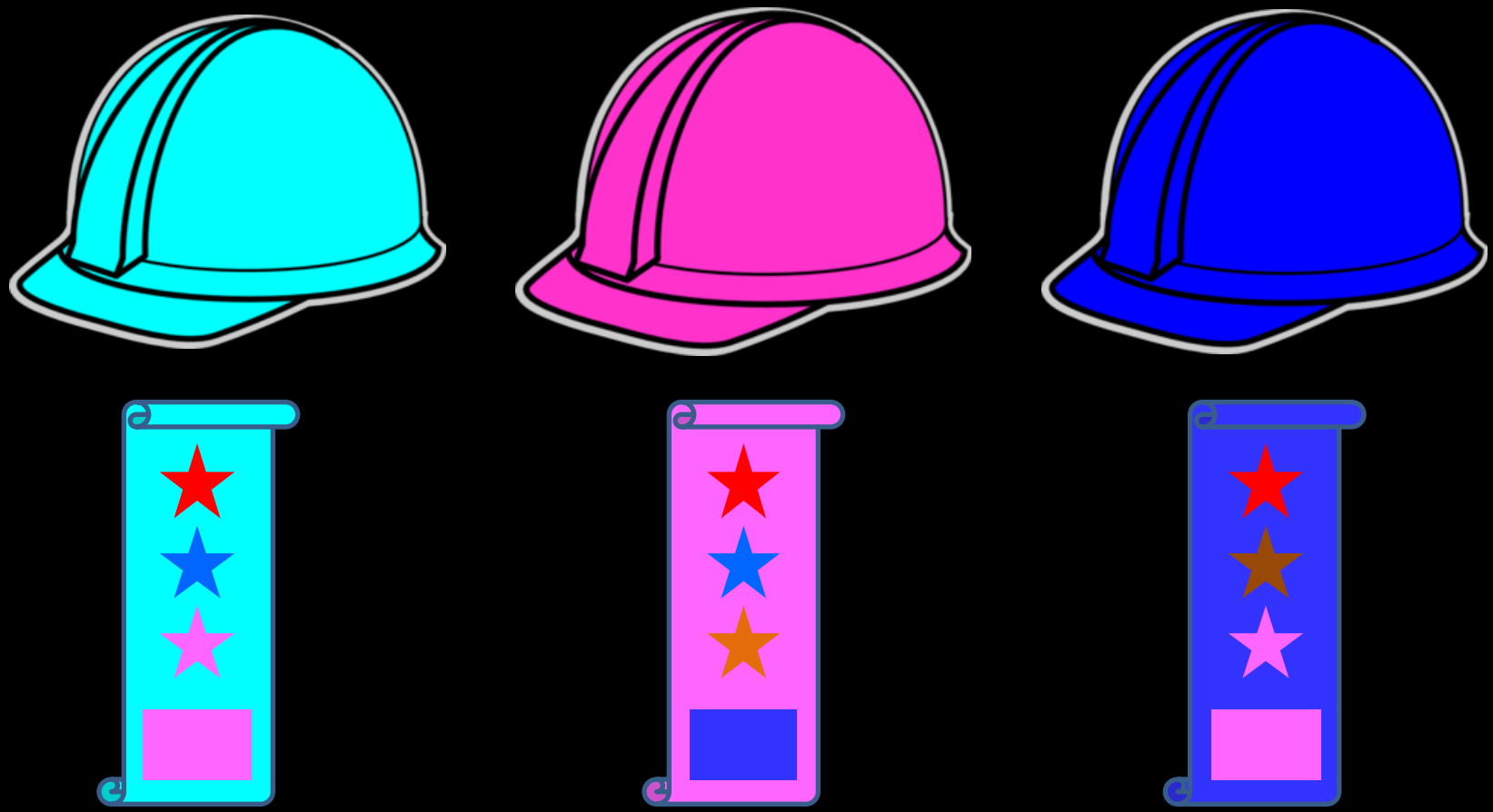


Block has contracts & new state

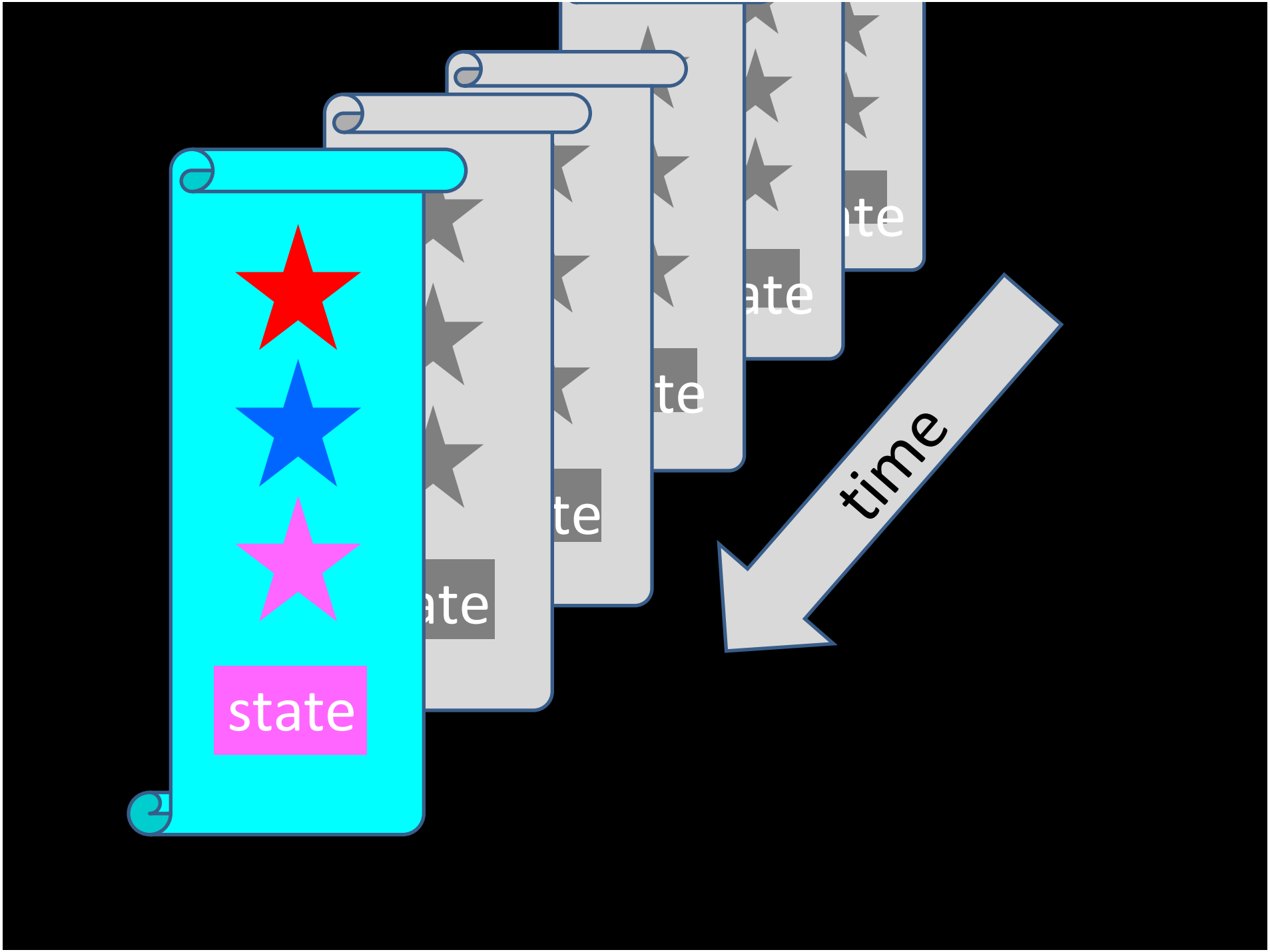




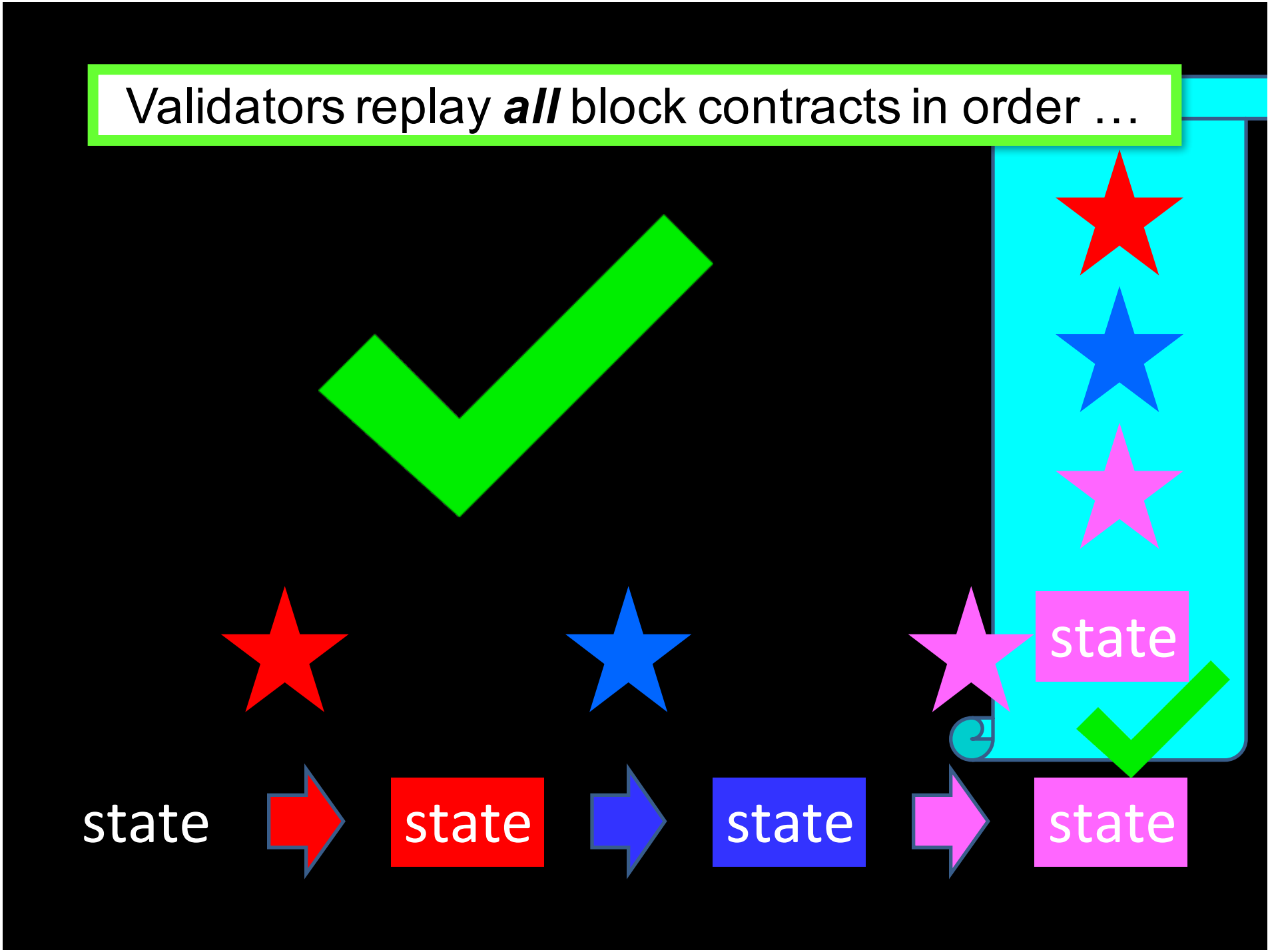
There can only be one…

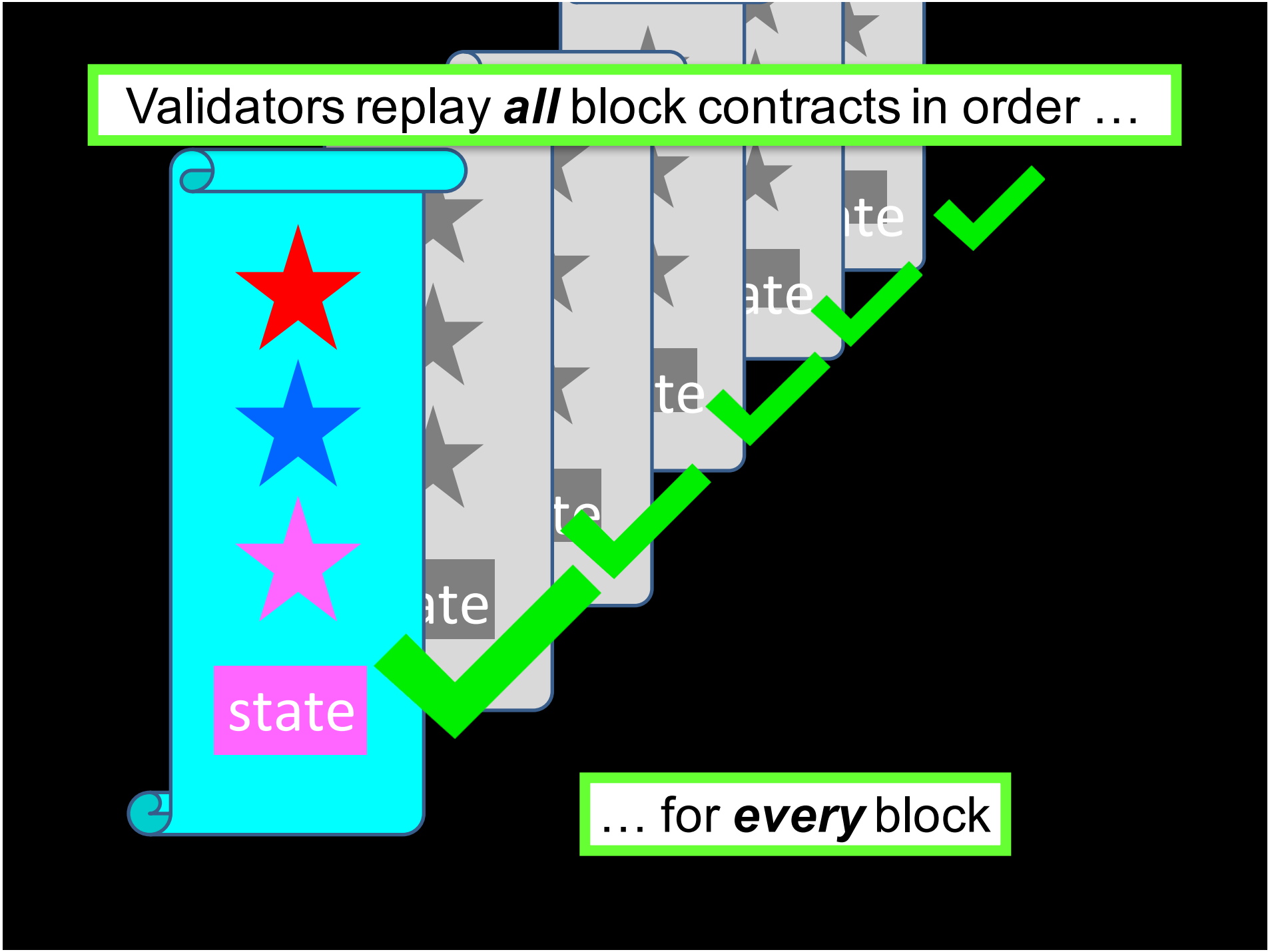


Miners compete to append *their* new block to the chain

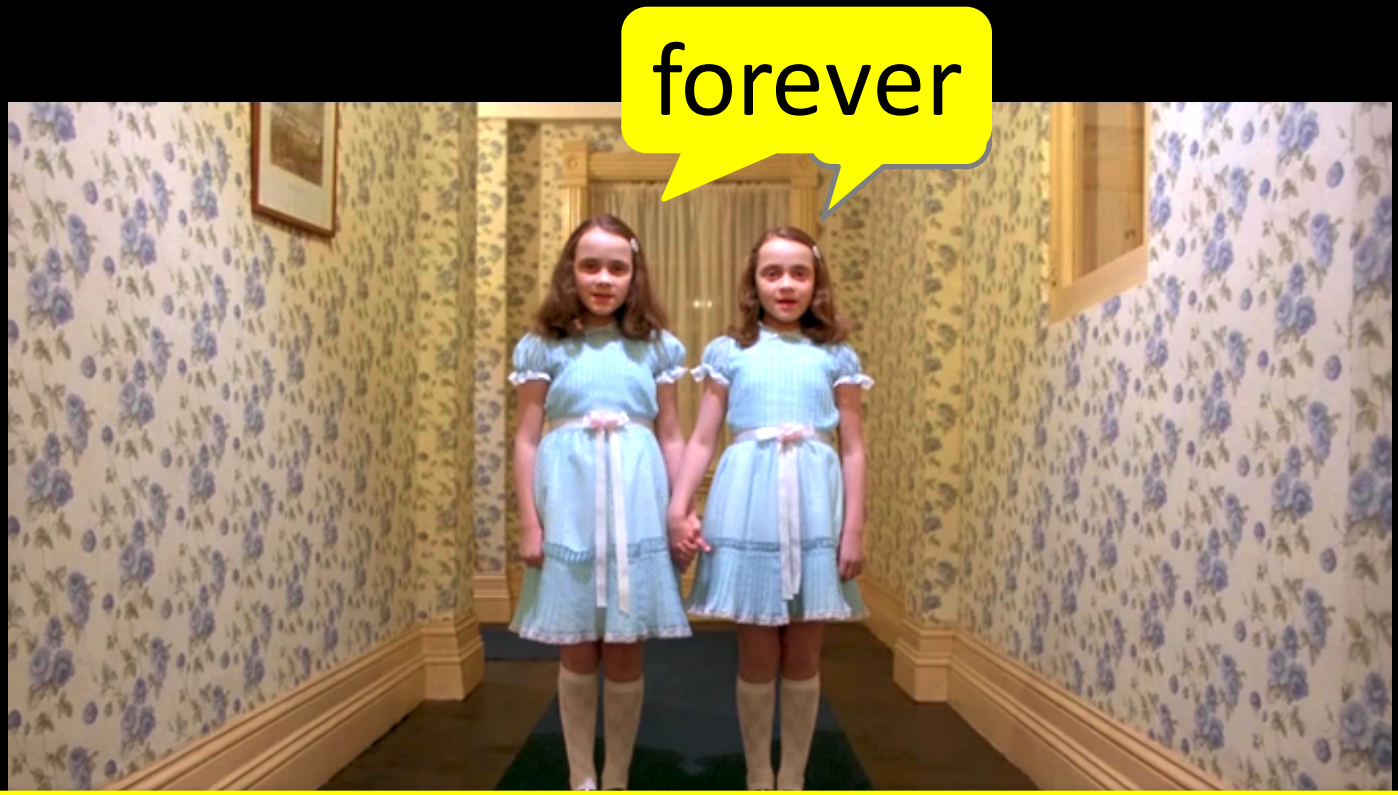

state
time

Validators replay ***all*** block contracts in order …



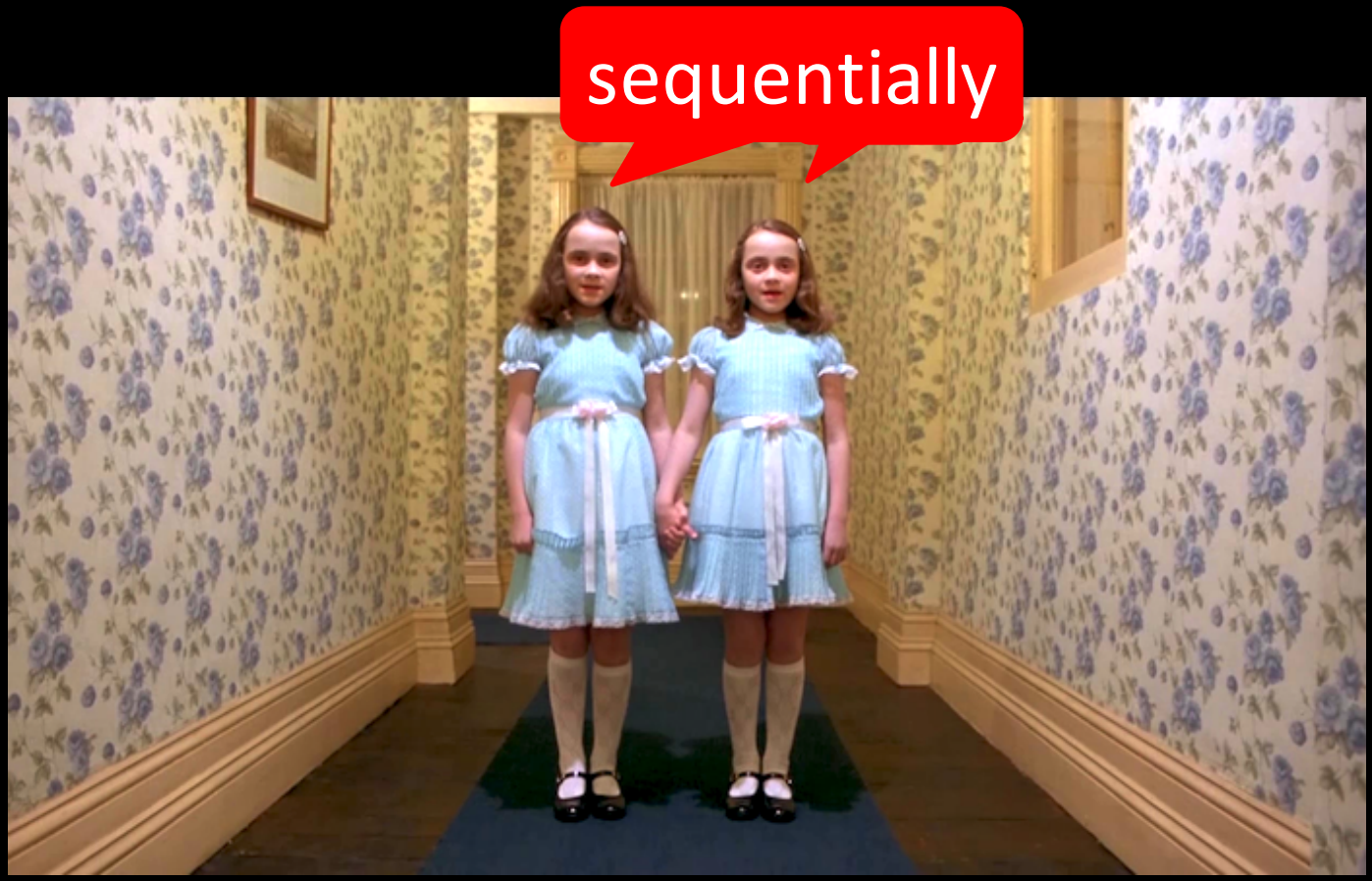

Validators replay *all* block contracts in order …
state
… for *every* block

Contracts are re-executed…

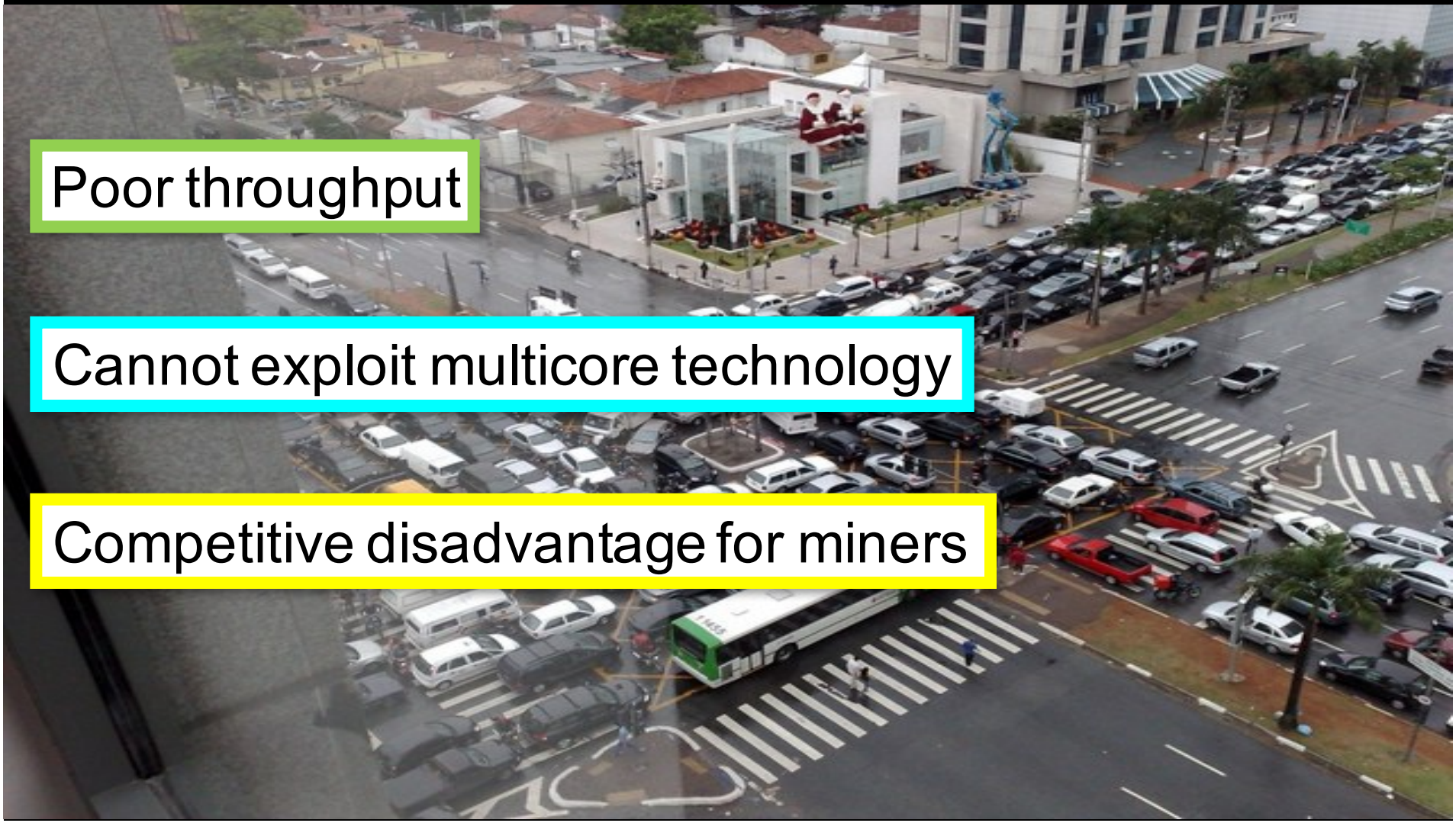


Every validator eventually executes *every* contract

Contracts are re-executed…



Why is sequential execution so wrong?



Poor throughput

Cannot exploit multicore technology

Competitive disadvantage for miners

Adding Concurrency

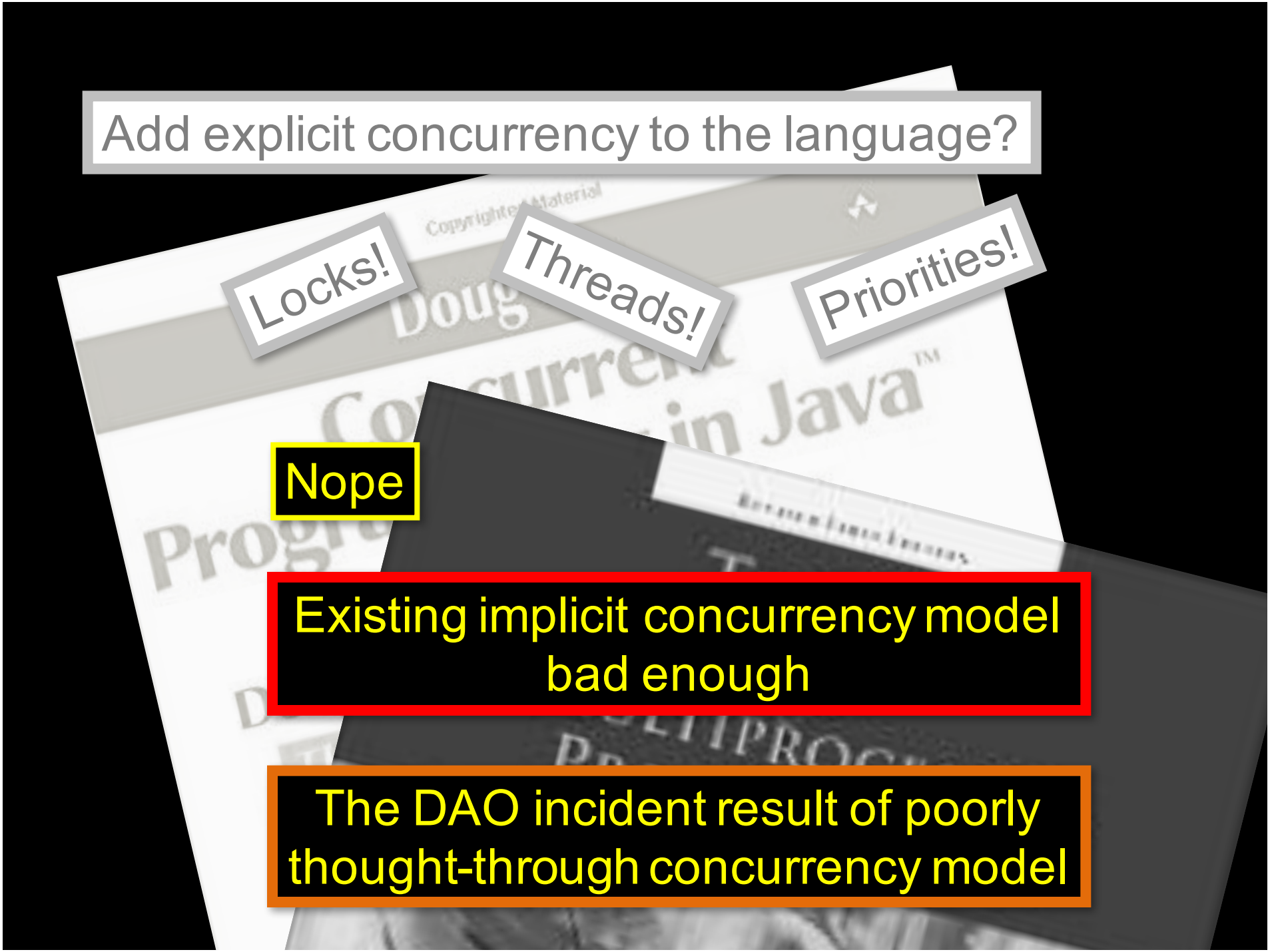
Naïve Concurrency?

Nope

Inconsistent shared state

Voters could vote twice

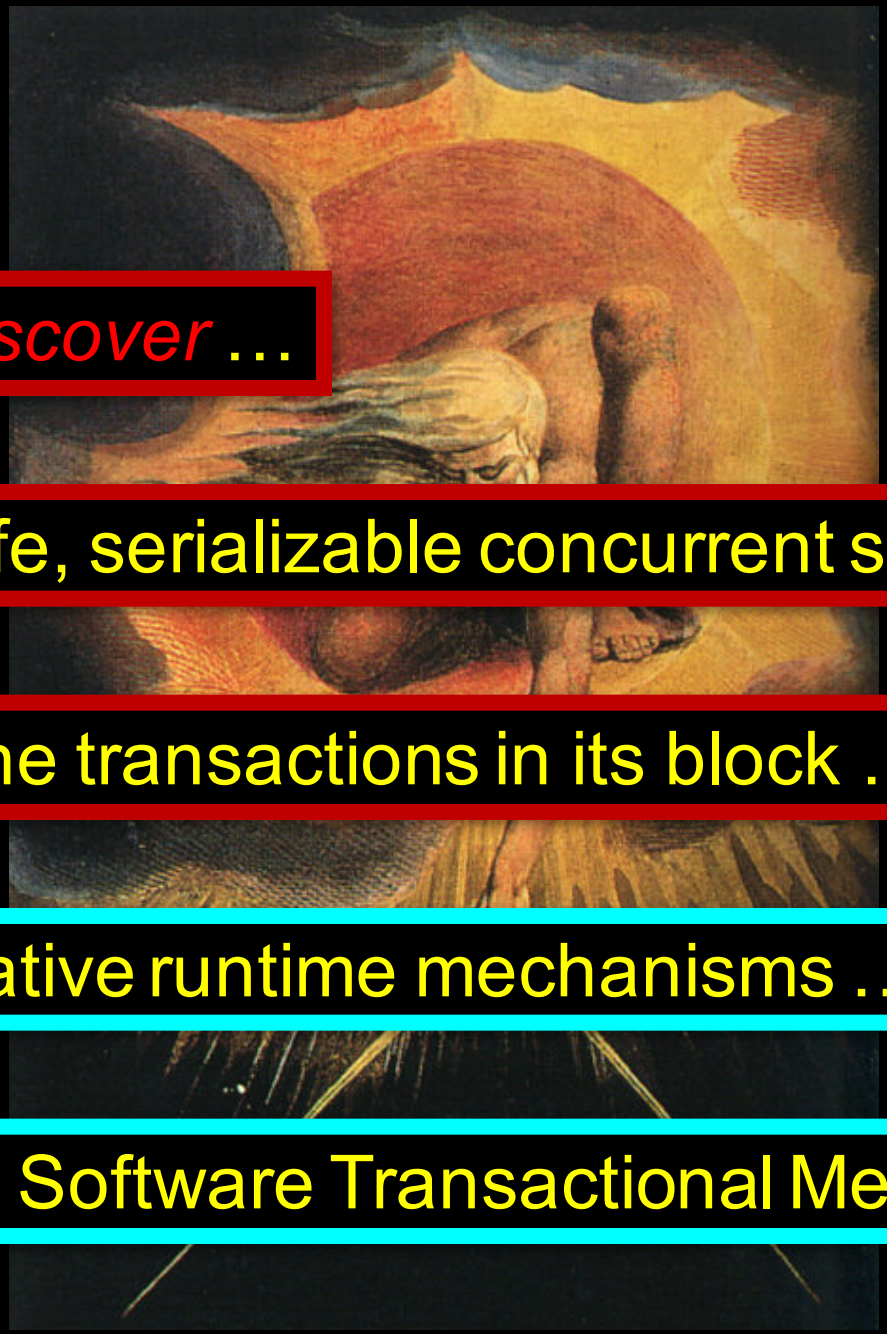

Add explicit concurrency to the language?
Locks!
Threads!
Priorities!
Doug Lea
Concurrent Programming in Java
Second Edition
Design Patterns
The Art of Multiprocessor Programming


Add explicit concurrency to the language?
Locks!
Threads!
Priorities!
Nope
Existing implicit concurrency model bad enough
The DAO incident result of poorly thought-through concurrency model


Concurrency via Static Analysis?
These contracts never conflict, so it's safe to run them concurrently


Concurrency via Static Analysis?
These contracts never conflict, so it's safe to run them concurrently
Nope
Undecidable in theory & Intractable in practice
Because contracts call other contracts …
Dynamic call graph …
Might have to inspect **every** contract **ever**

Big Idea #1



Let miners *discover* …

a safe, serializable concurrent schedule …

for the transactions in its block …

using speculative runtime mechanisms …

adapted from Software Transactional Memory.

Instrument shared objects & variables

E.g., locks on methods and accessors

Function are atomic sections

Conflict detected?

Delay or restart one thread

Keep track of “happens before”

Result is safe concurrent schedule + description

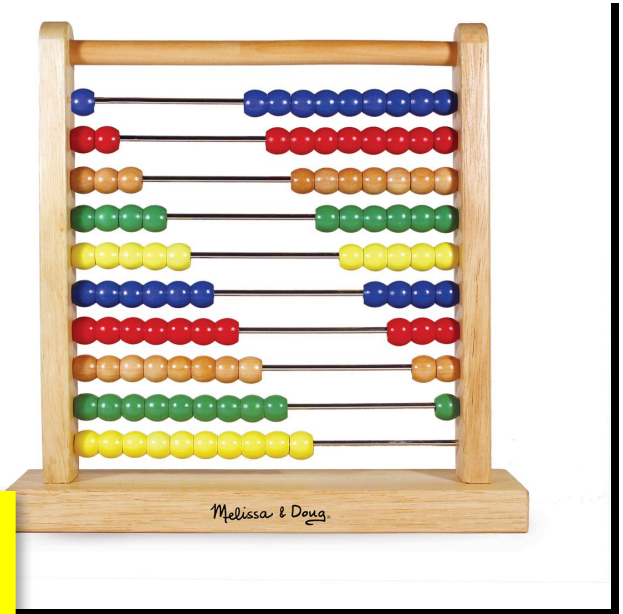
Positive

Usually, conflict is rare

Easy concurrent executions

Less delay is competitive advantage

Better HW usage, less energy, etc.



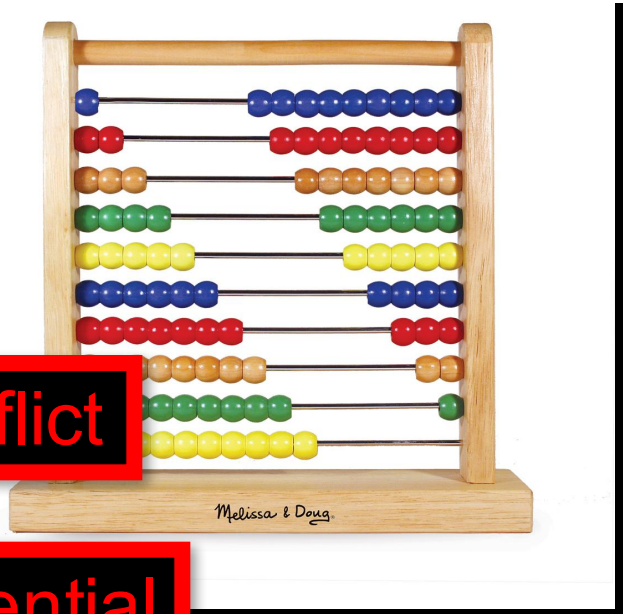
Negative

Sometimes transactions do conflict

Executions must be sequential

Synchronization overhead means delay

But here are many tricks …



Positive

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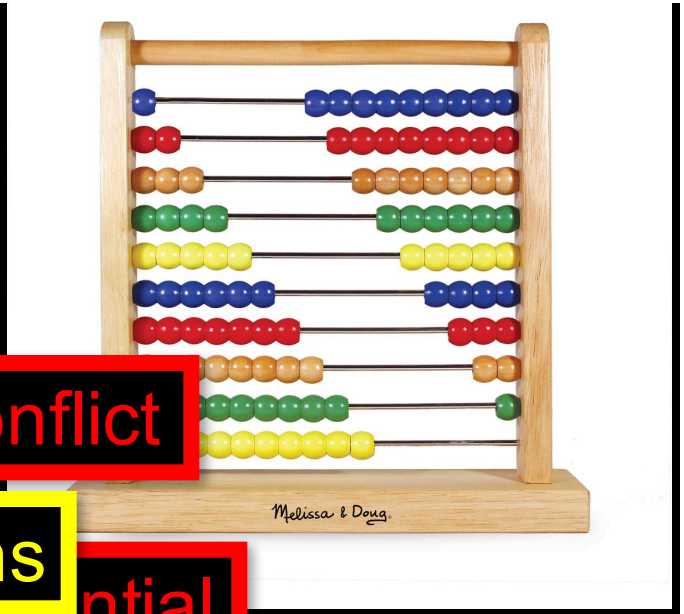


Take your choice

What about validators

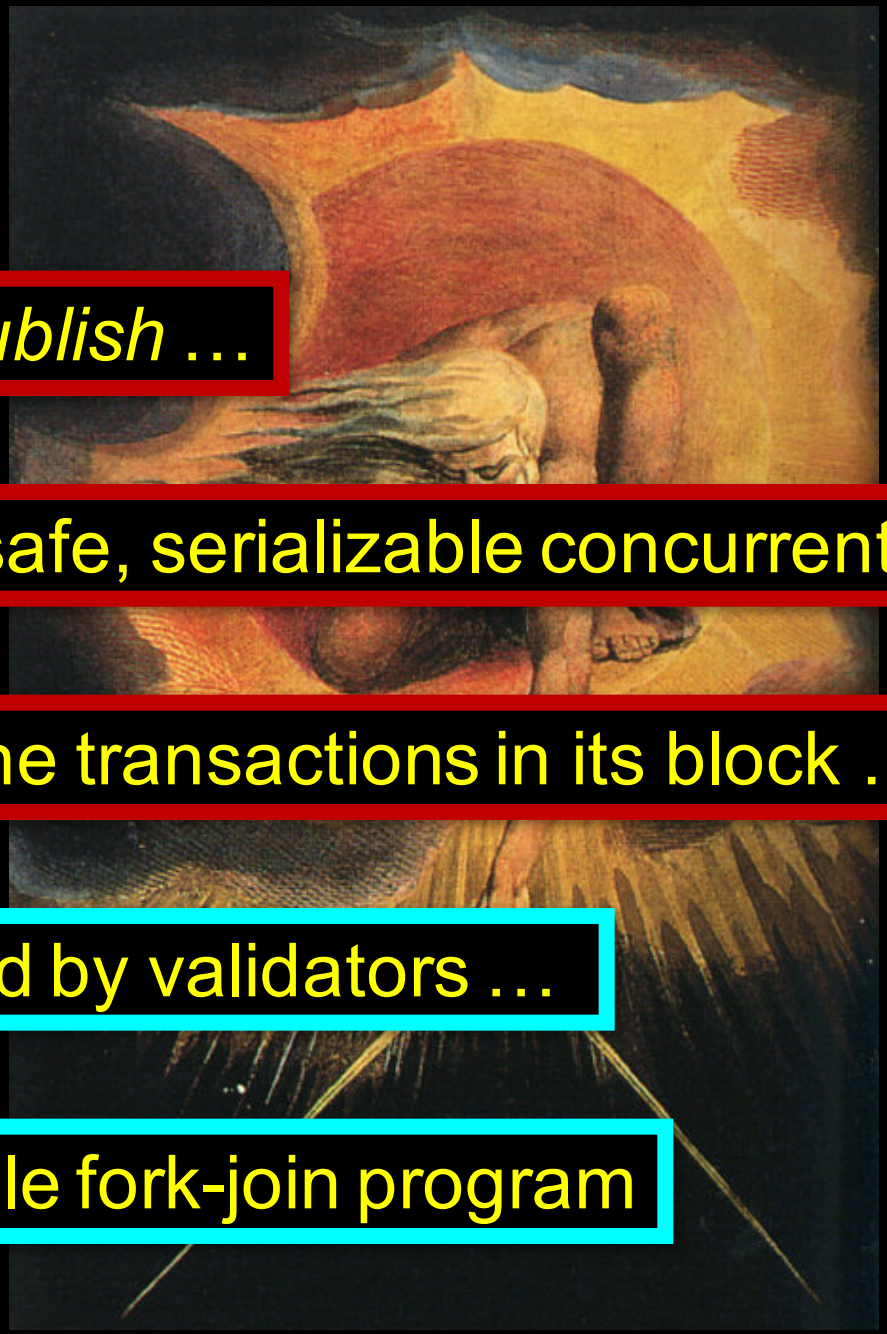
Cannot mimic miners by discovering schedules

Parallel executions non-deterministic

Might find a different safe concurrent schedule

Or resort to sequential execution

Big Idea #2



Let miners *publish* …

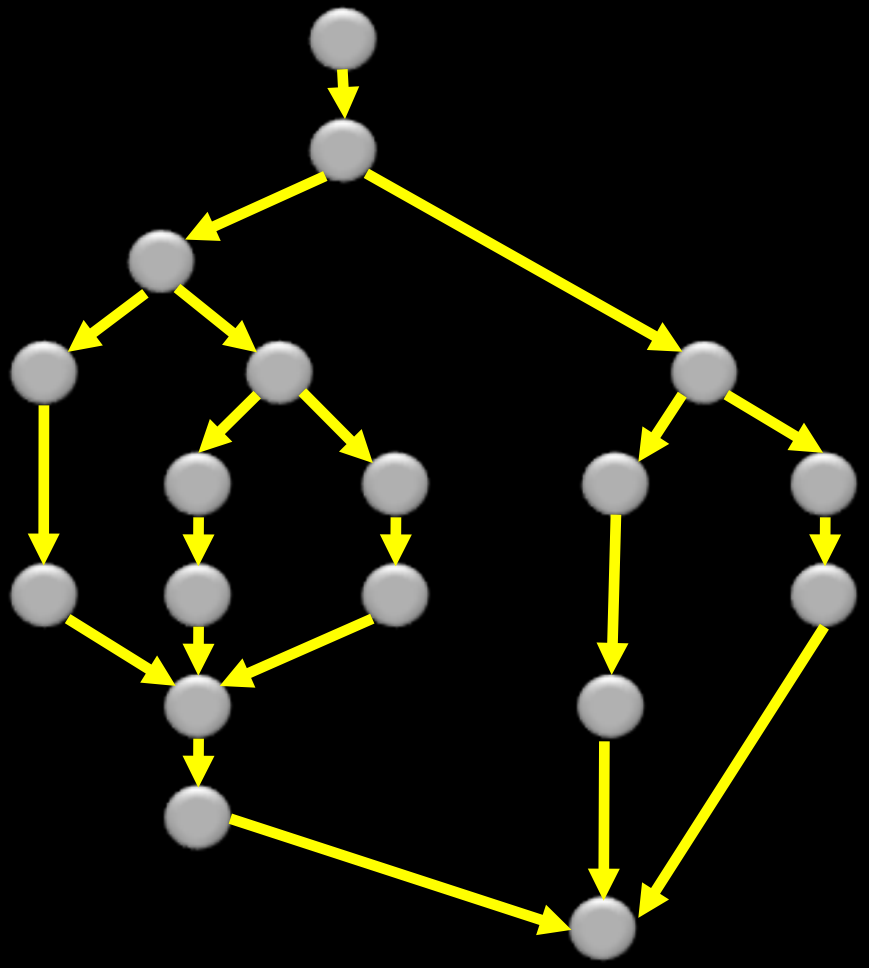
the safe, serializable concurrent schedule …

for the transactions in its block …

to be replayed by validators …

as a checkable fork-join program

Generate a Fork-Join Program



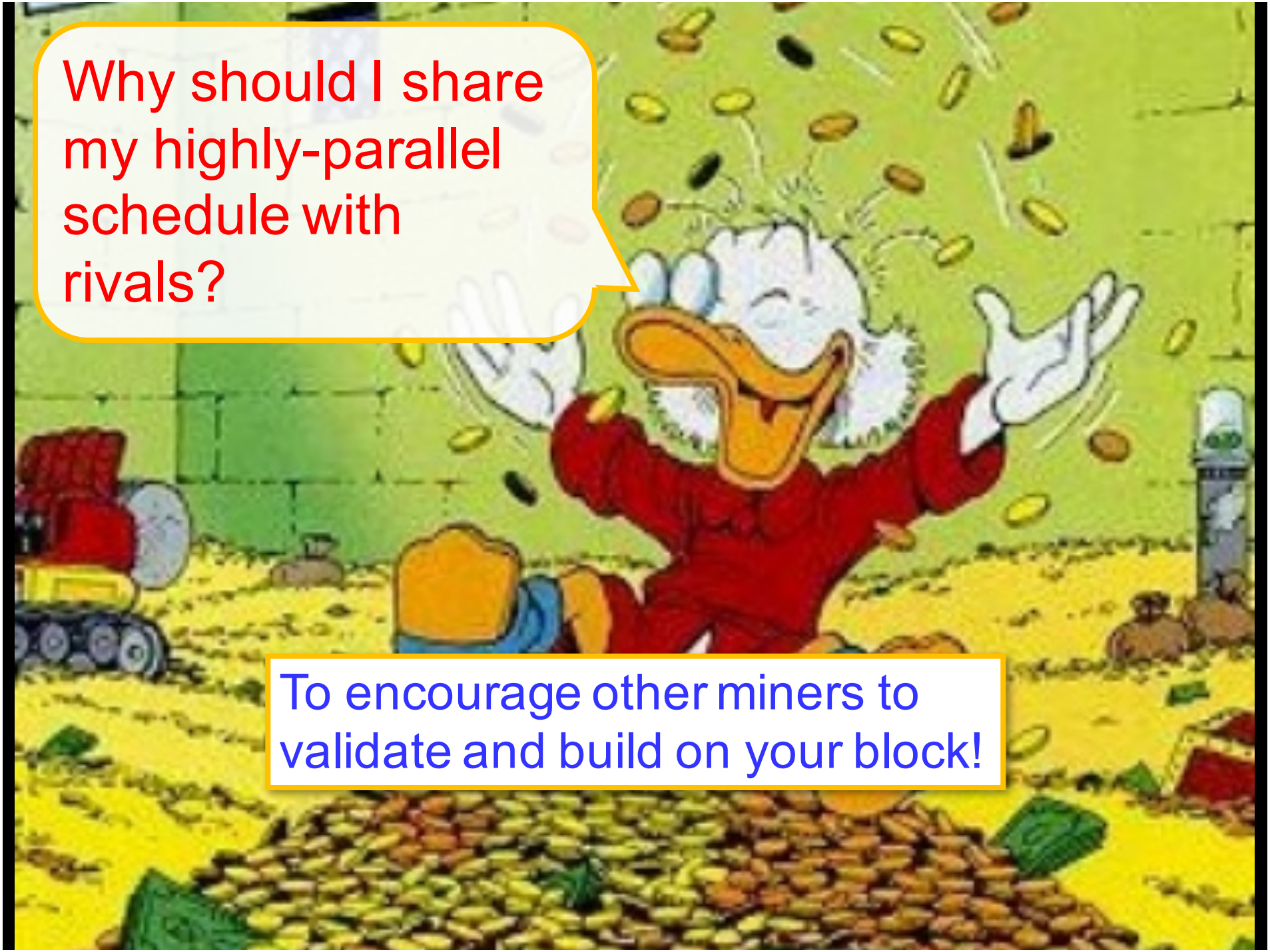
Similar to CILK model

Efficient work-stealing scheduler

Can check validity

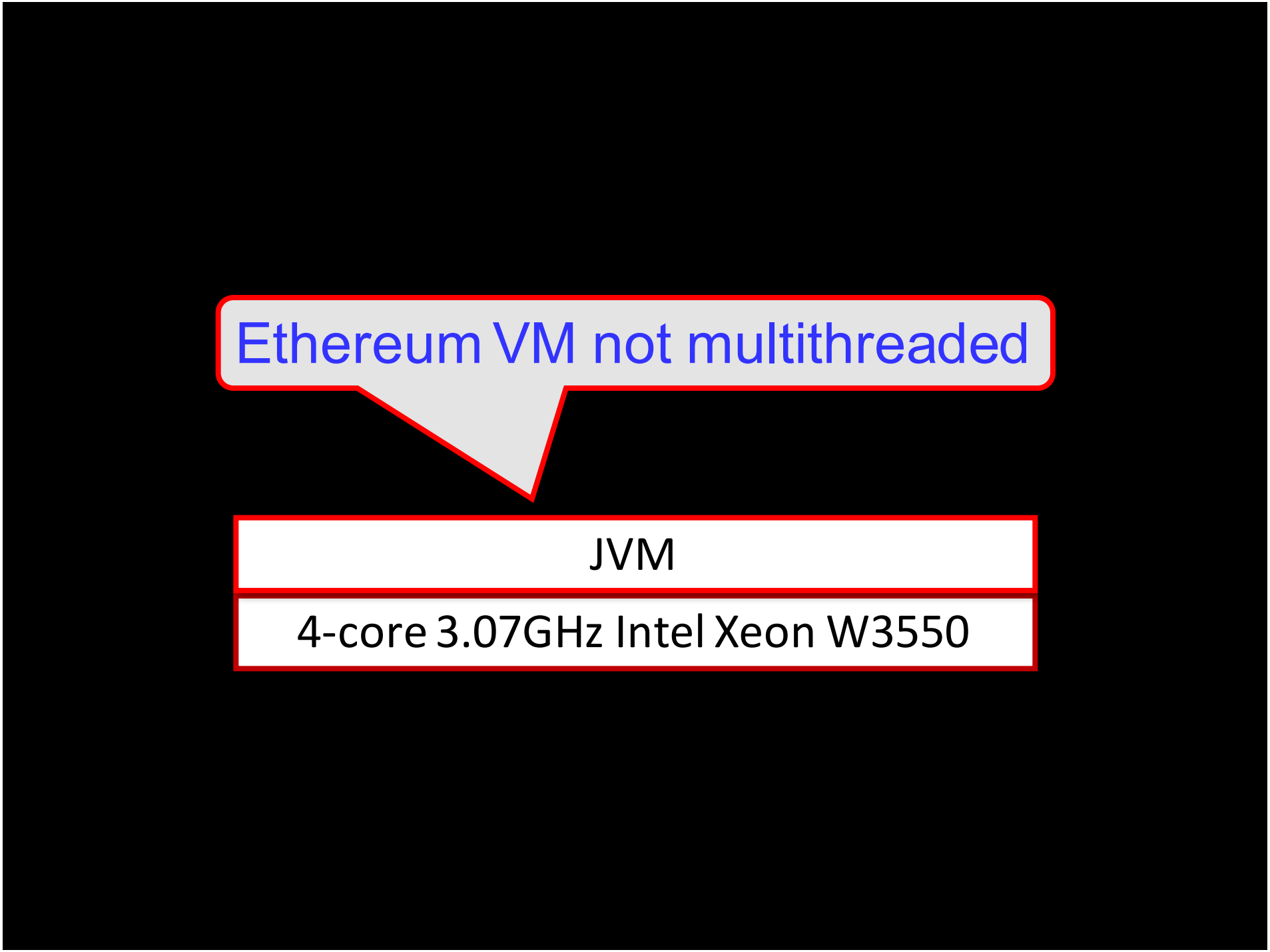
No locks, rollbacks

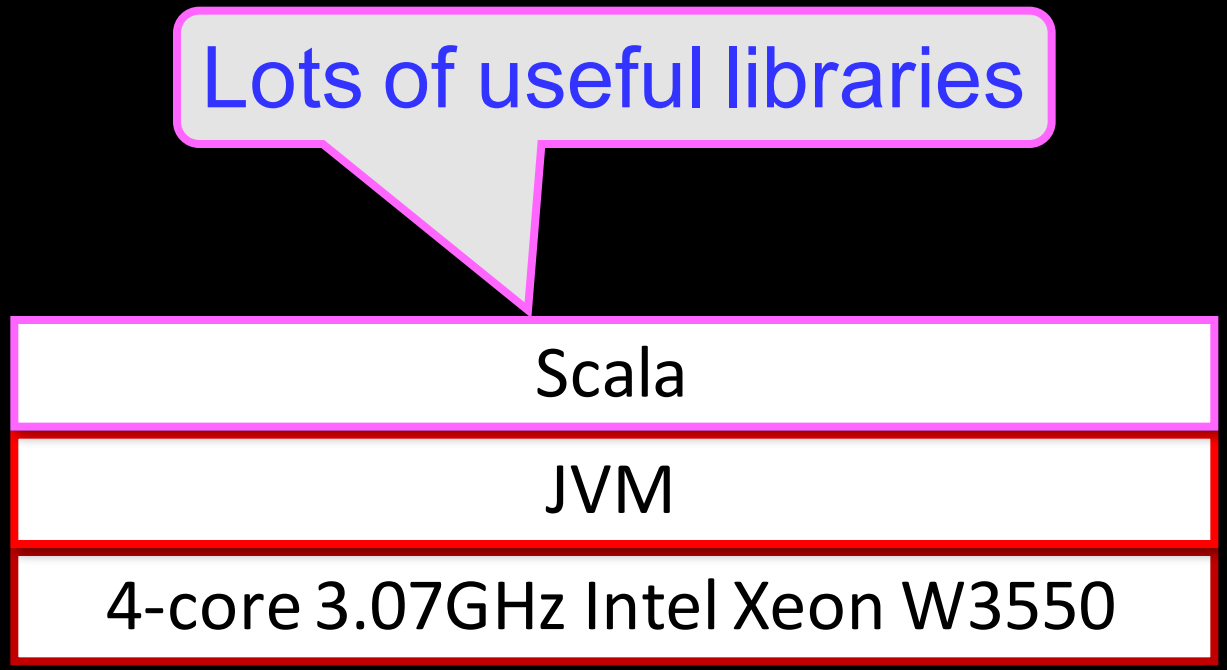
deterministic

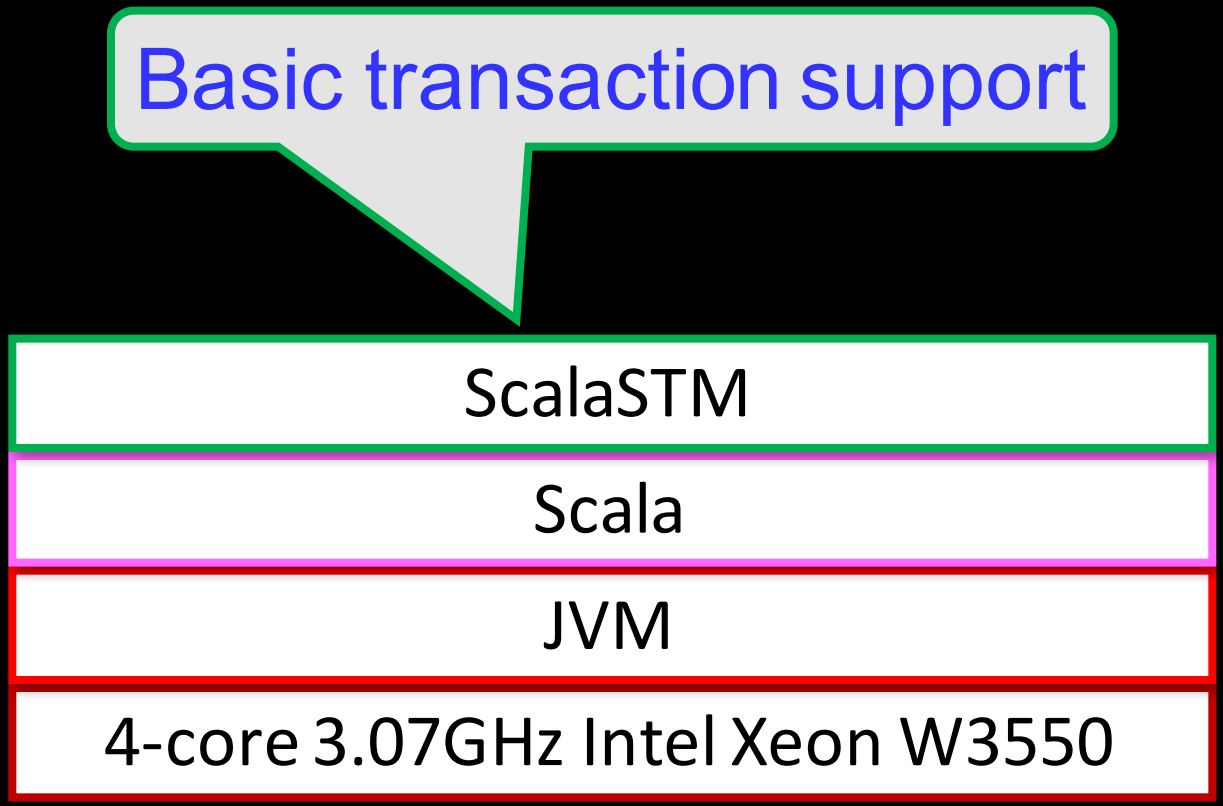

Why should I share my highly-parallel schedule with rivals?
To encourage other miners to validate and build on your block!

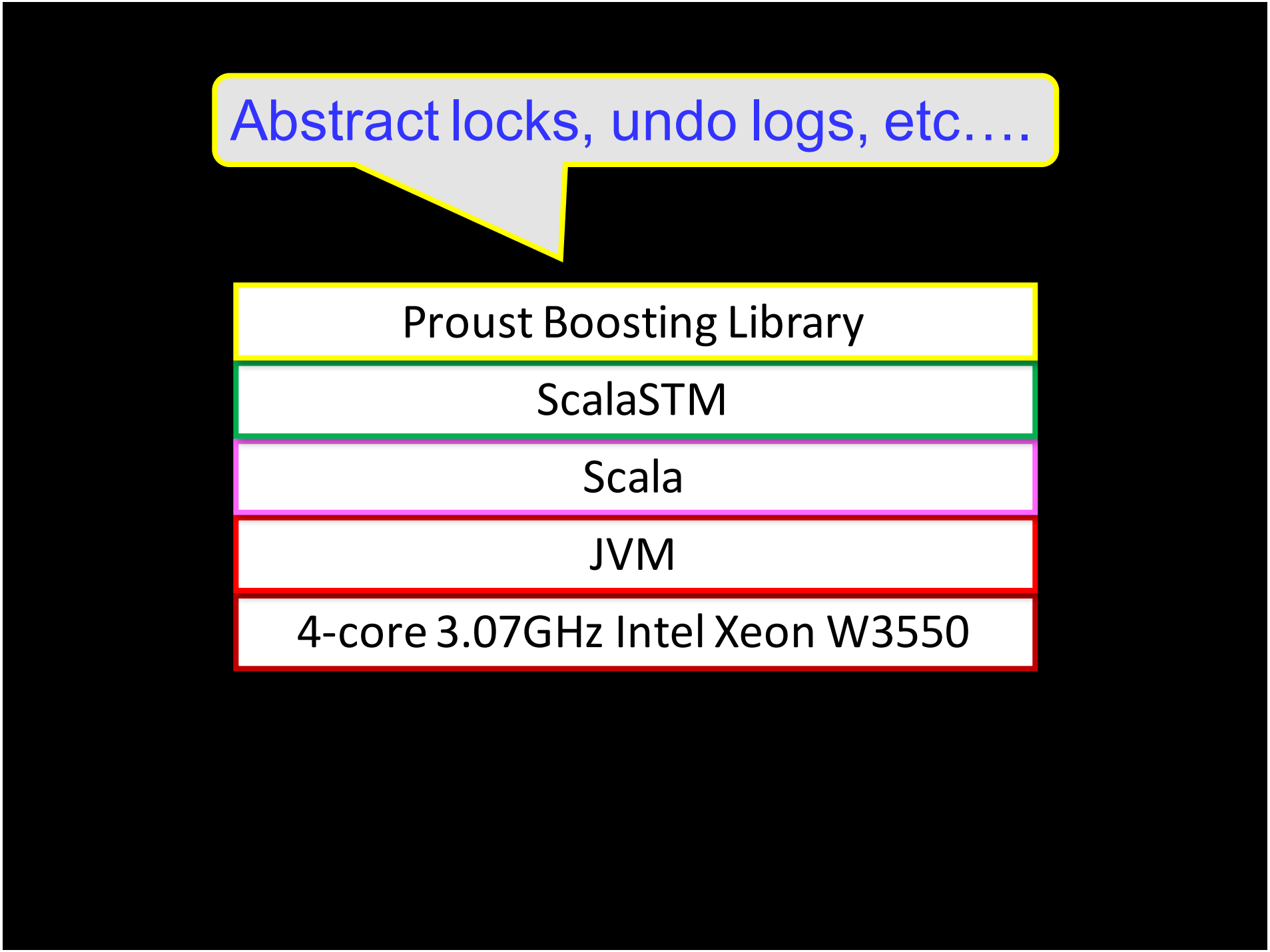
Prototype and Evaluation


Available hardware
4-core 3.07GHz Intel Xeon W3550


Ethereum VM not multithreaded
JVM
4-core 3.07GHz Intel Xeon W3550


Lots of useful libraries
Scala
JVM
4-core 3.07GHz Intel Xeon W3550


Basic transaction support
ScalaSTM
Scala
JVM
4-core 3.07GHz Intel Xeon W3550


Abstract locks, undo logs, etc.…
Proust Boosting Library
ScalaSTM
Scala
JVM
4-core 3.07GHz Intel Xeon W3550

Benchmarks

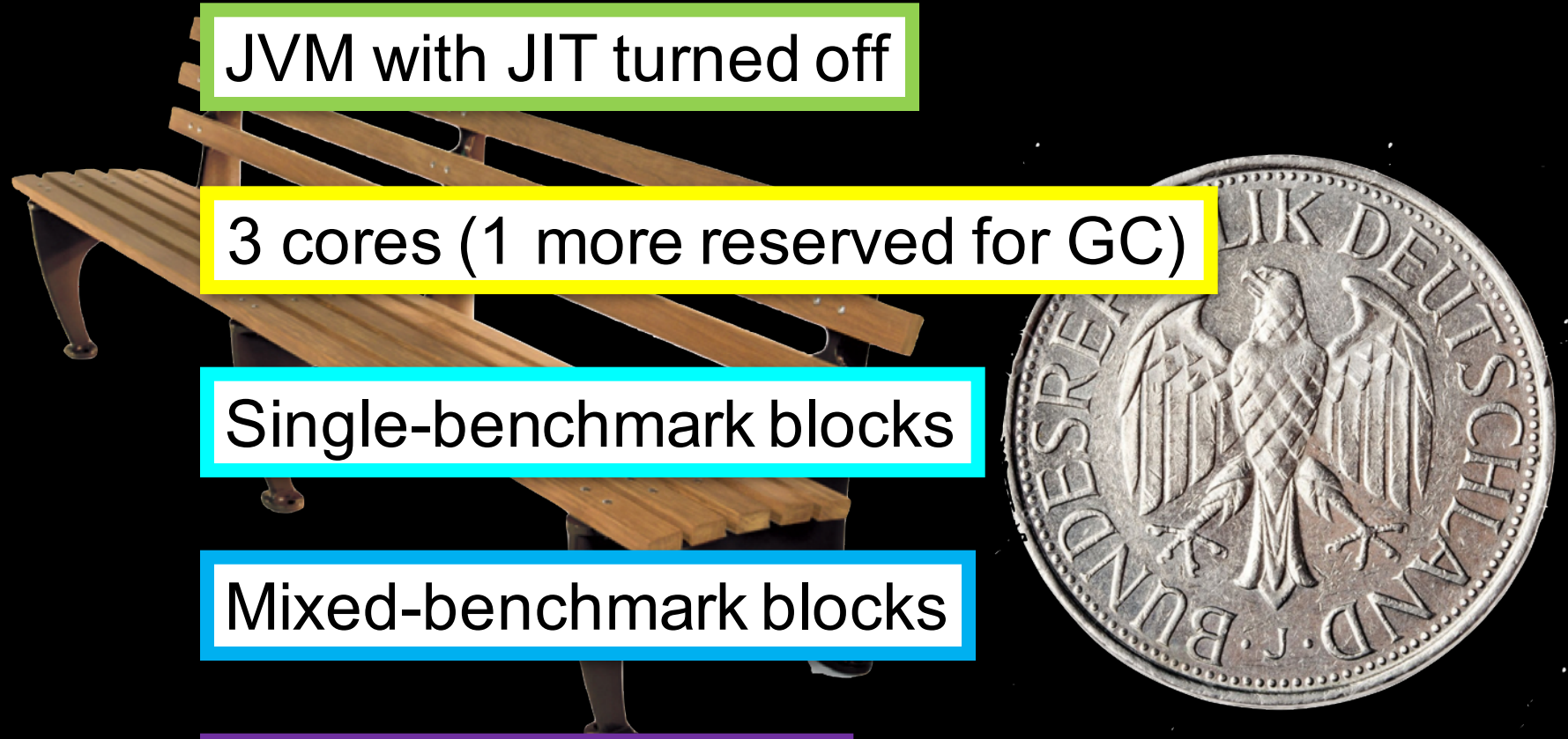
JVM with JIT turned off

3 cores (1 more reserved for GC)

Single-benchmark blocks

Mixed-benchmark blocks

Tunable Conflict rate



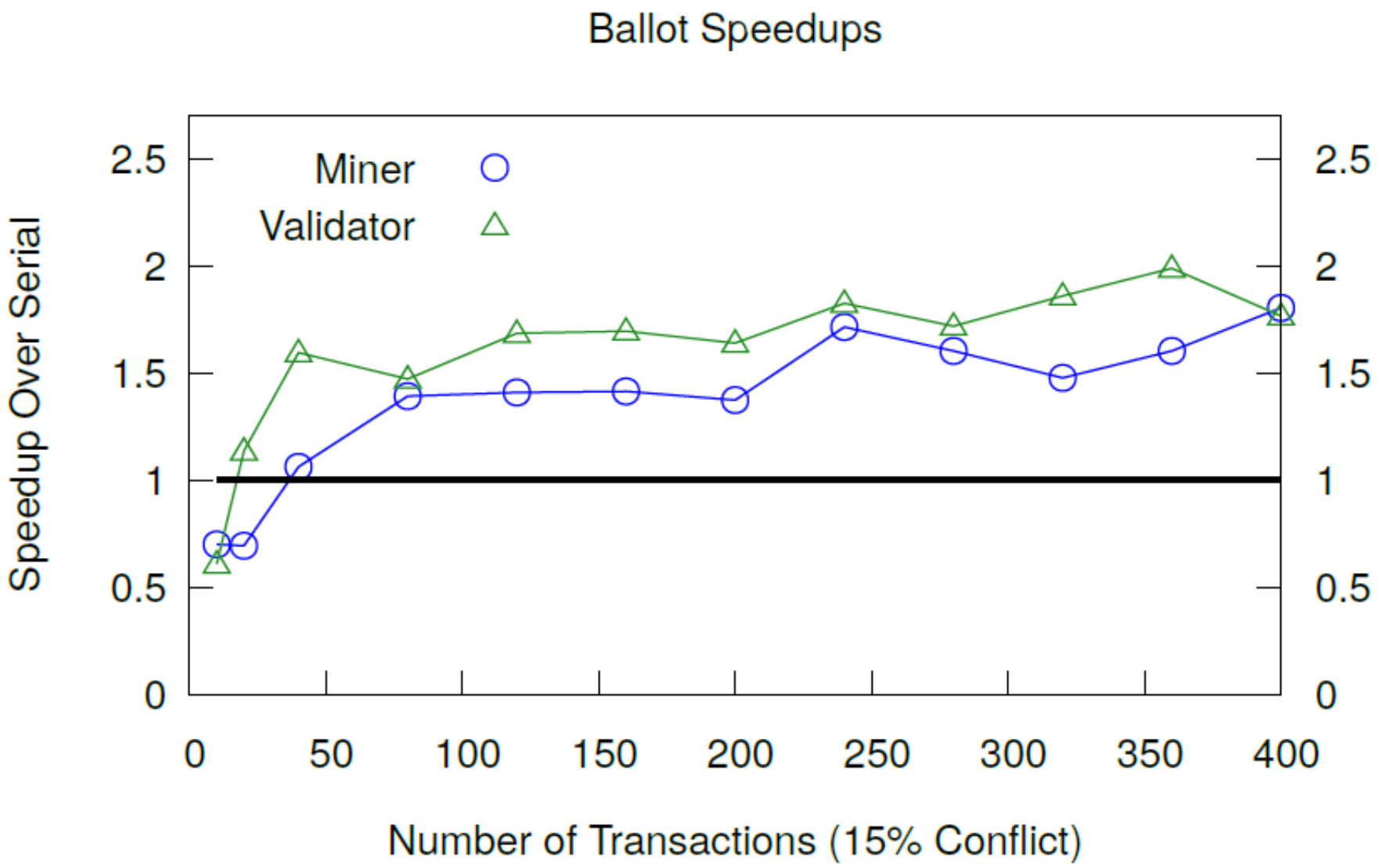
Benchmark #1: Ballot

From Solidity documentation

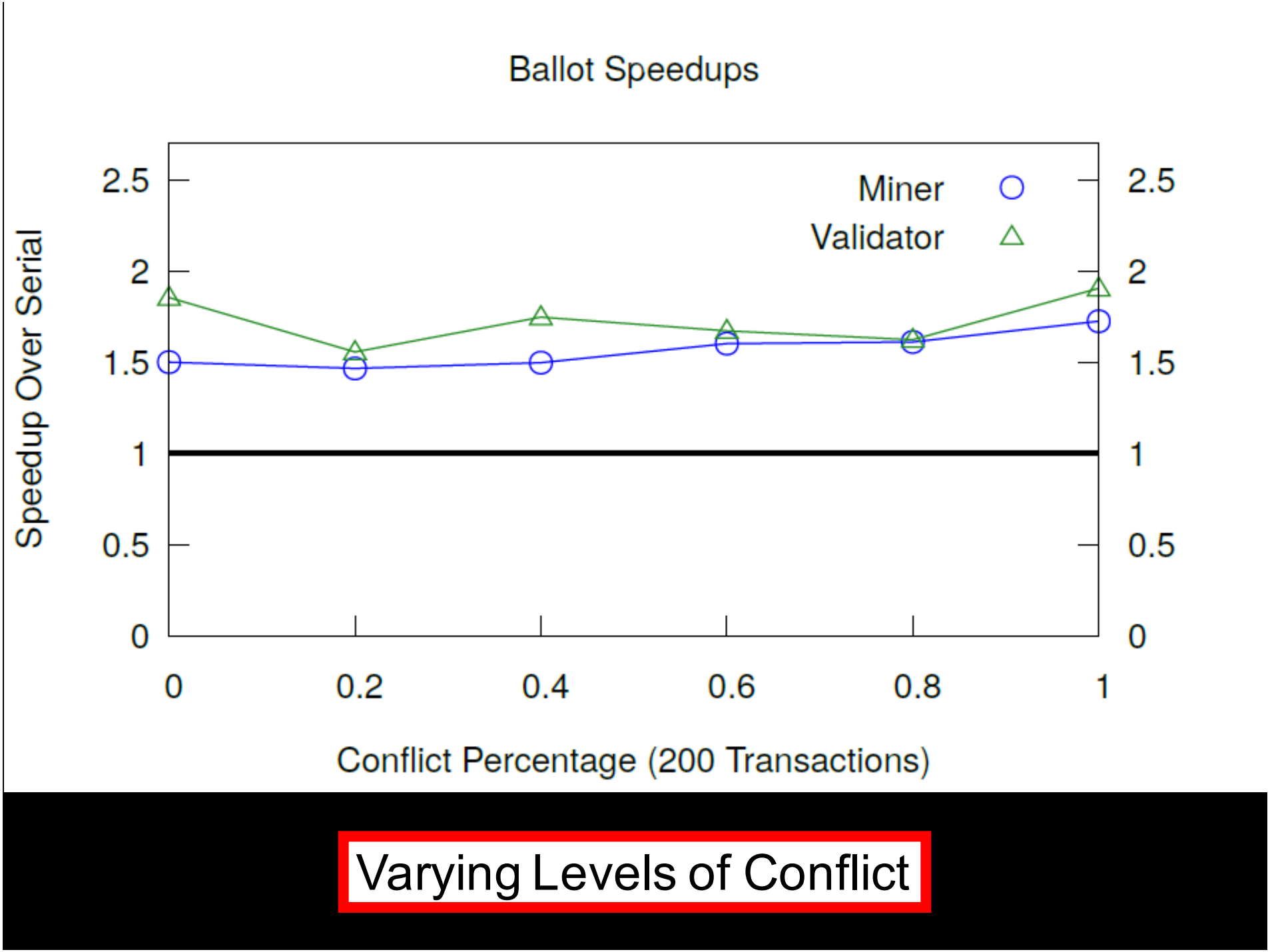
Benchmark: all voters registered, vote only

Shared state: voter mapping

Tunable Conflict = double voting



Varying Transactions per Block



Varying Levels of Conflict

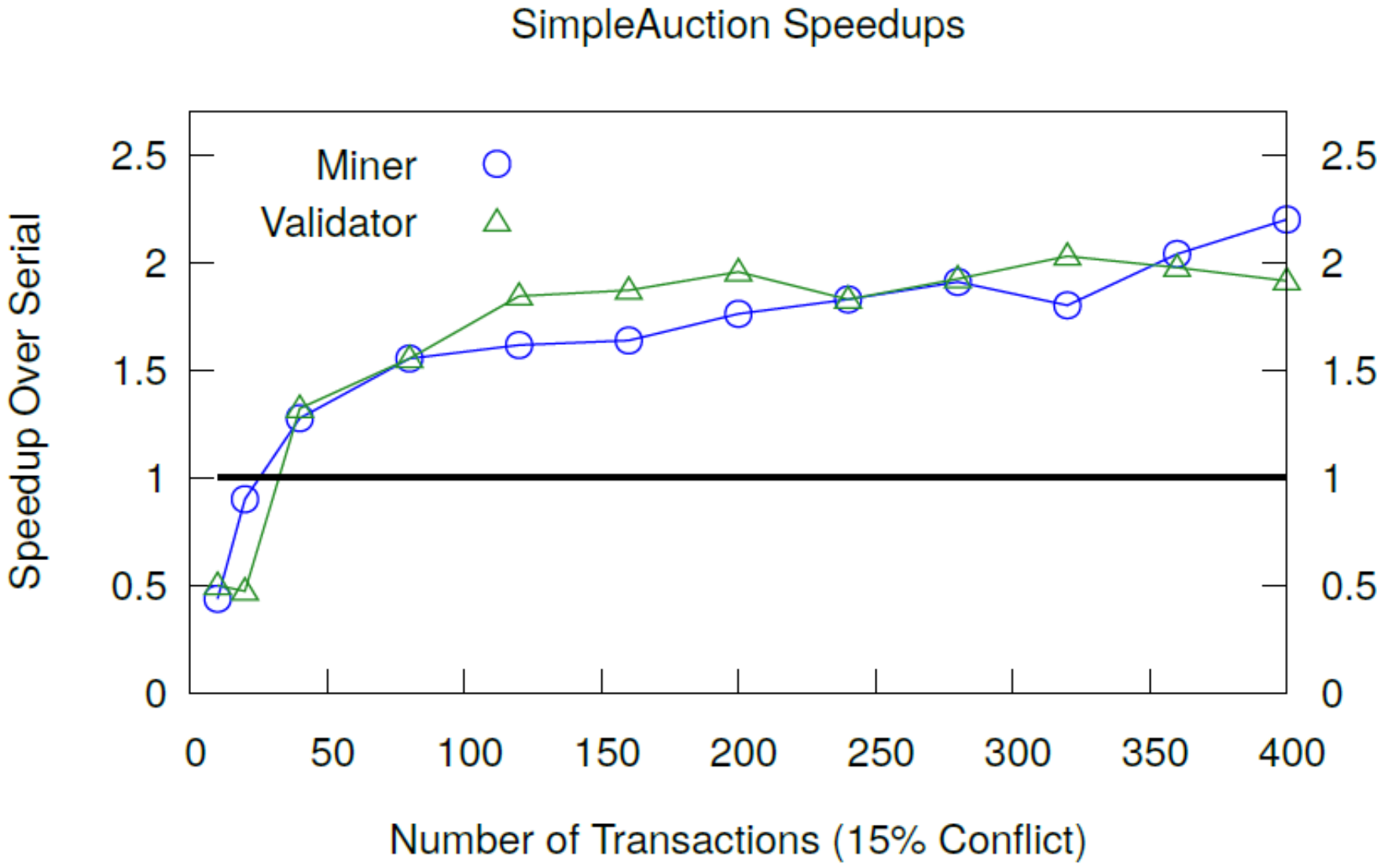
Benchmark #2: SimpleAuction

From Solidity documentation

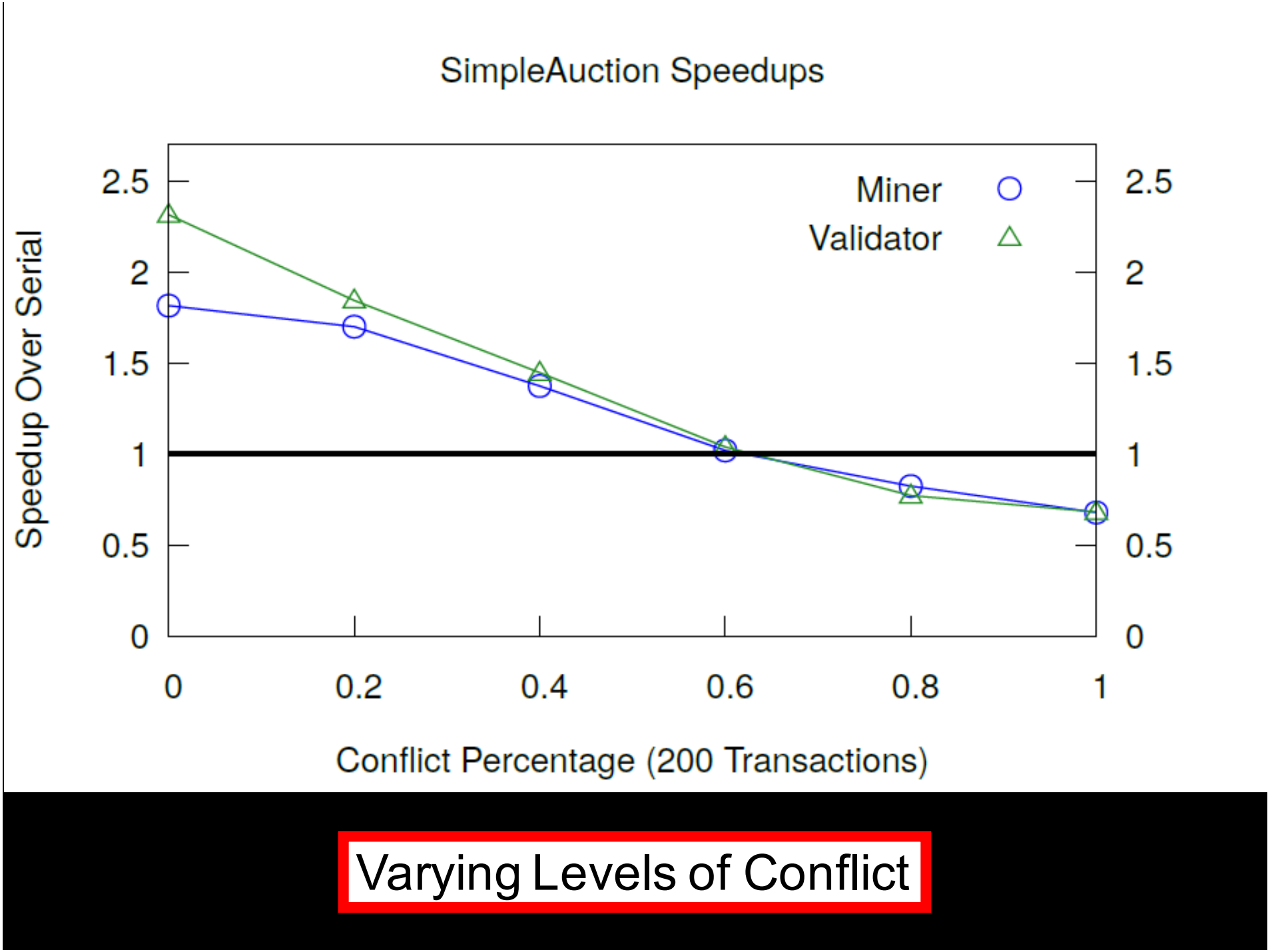
Benchmark: bidders bid, request refunds

Shared state: maxBid

Tunable Conflict = bidPlusOne() vs refund



Varying Transactions per Block



Varying Levels of Conflict

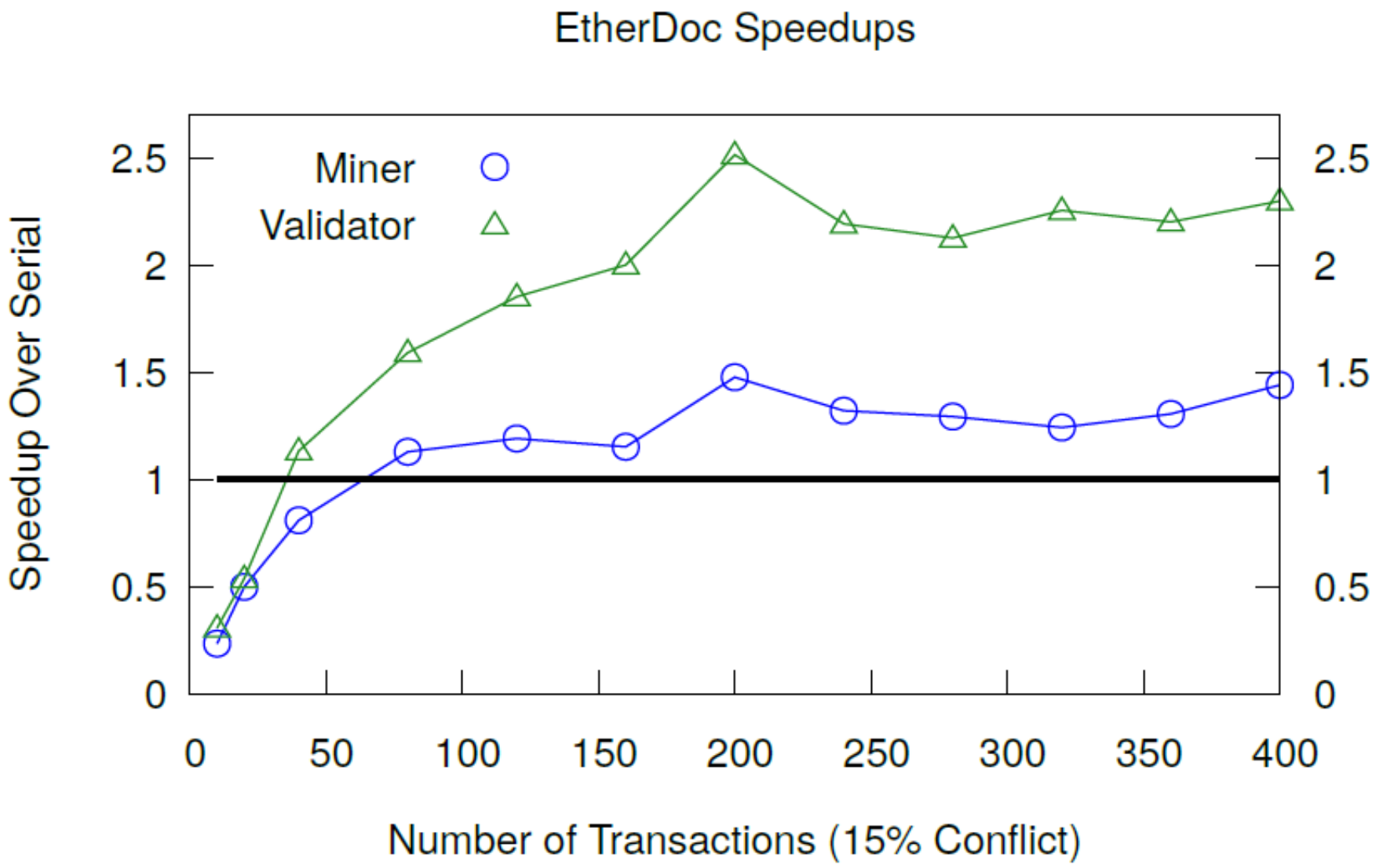
Benchmark #3: EtherDoc

From website

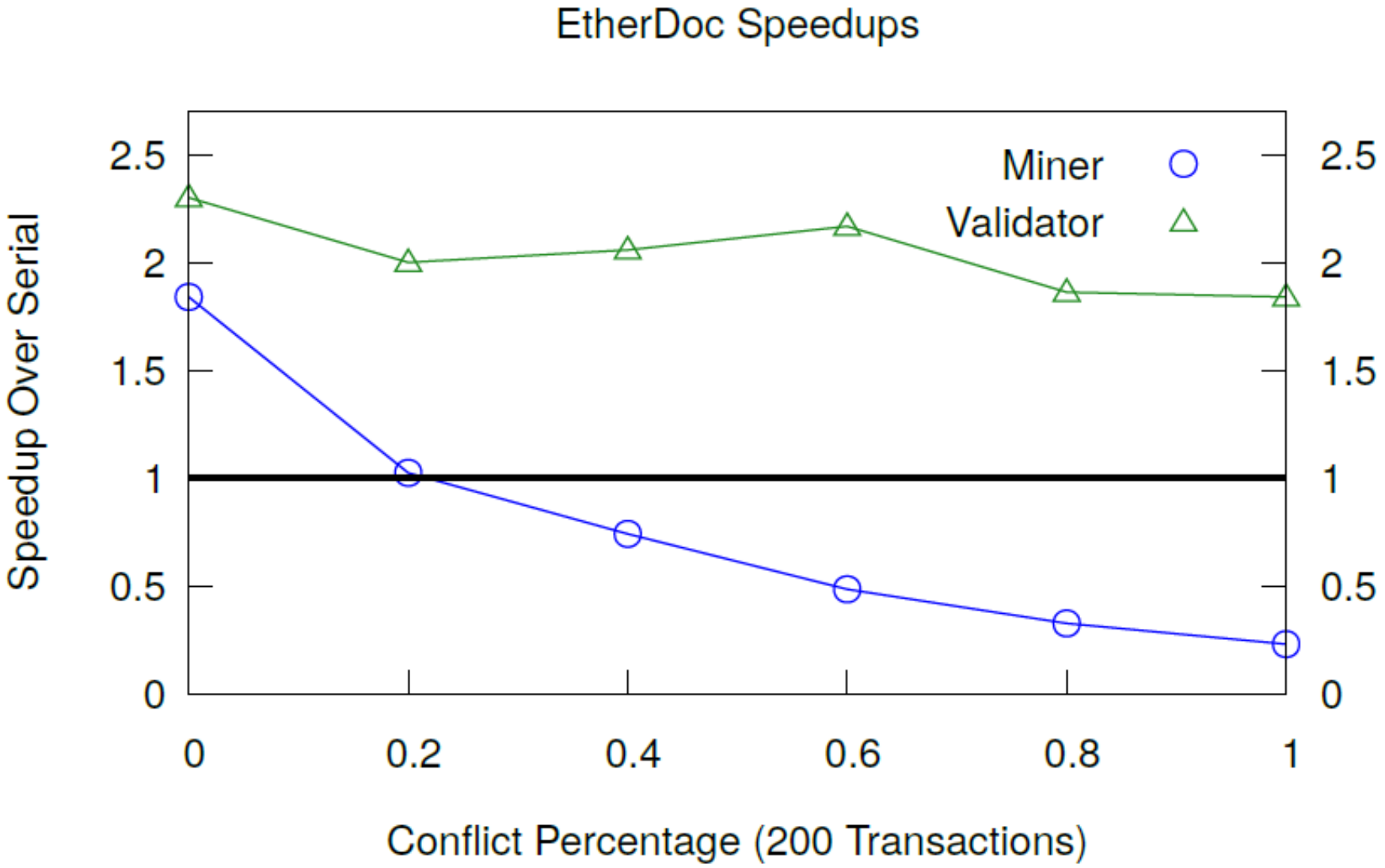
Tracks Document Metadata (including owner)

Shared state: owner's list of docs

Tunable Conflict = transfer vs query



Varying Transactions per Block

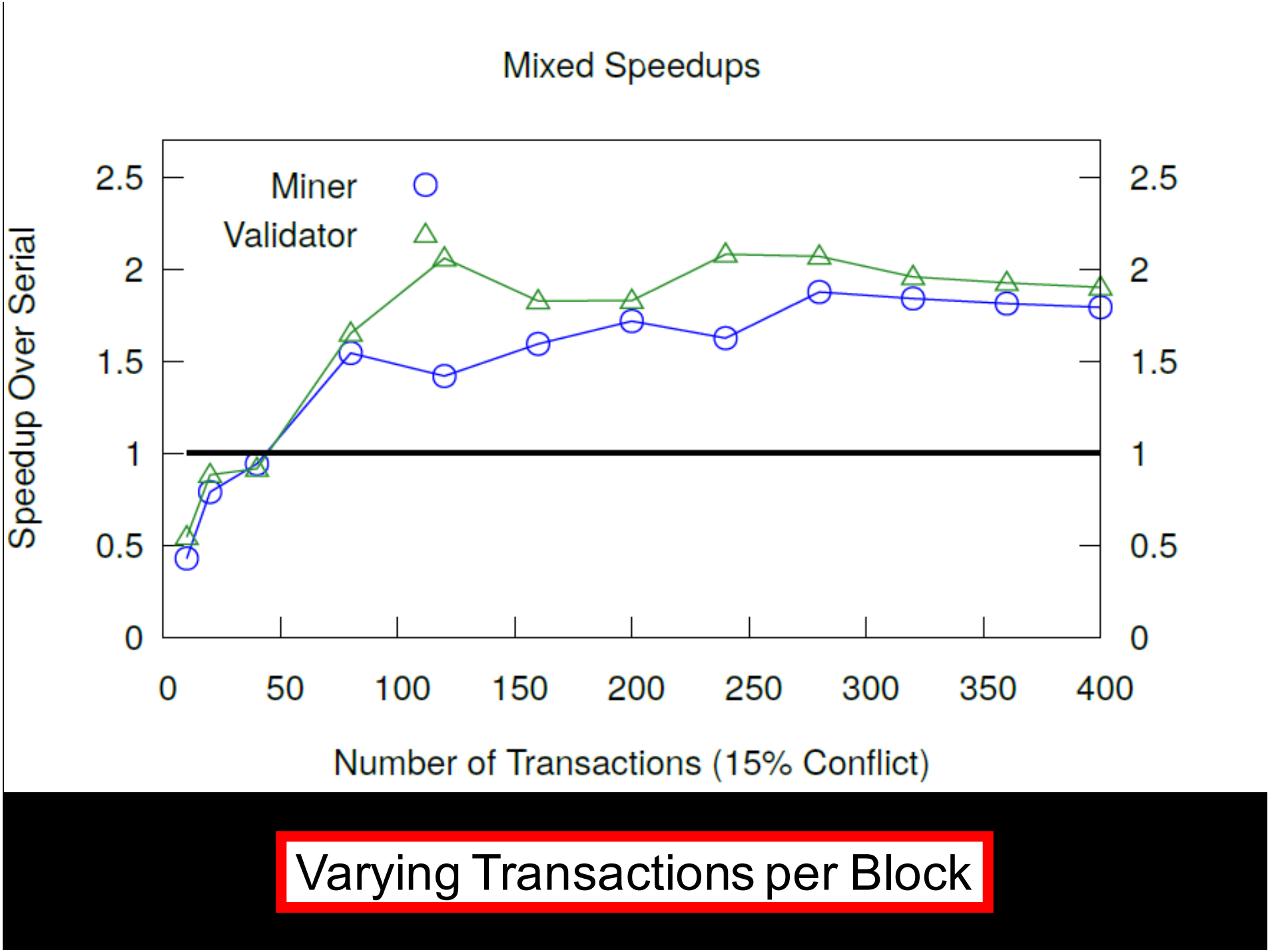


Varying Levels of Conflict

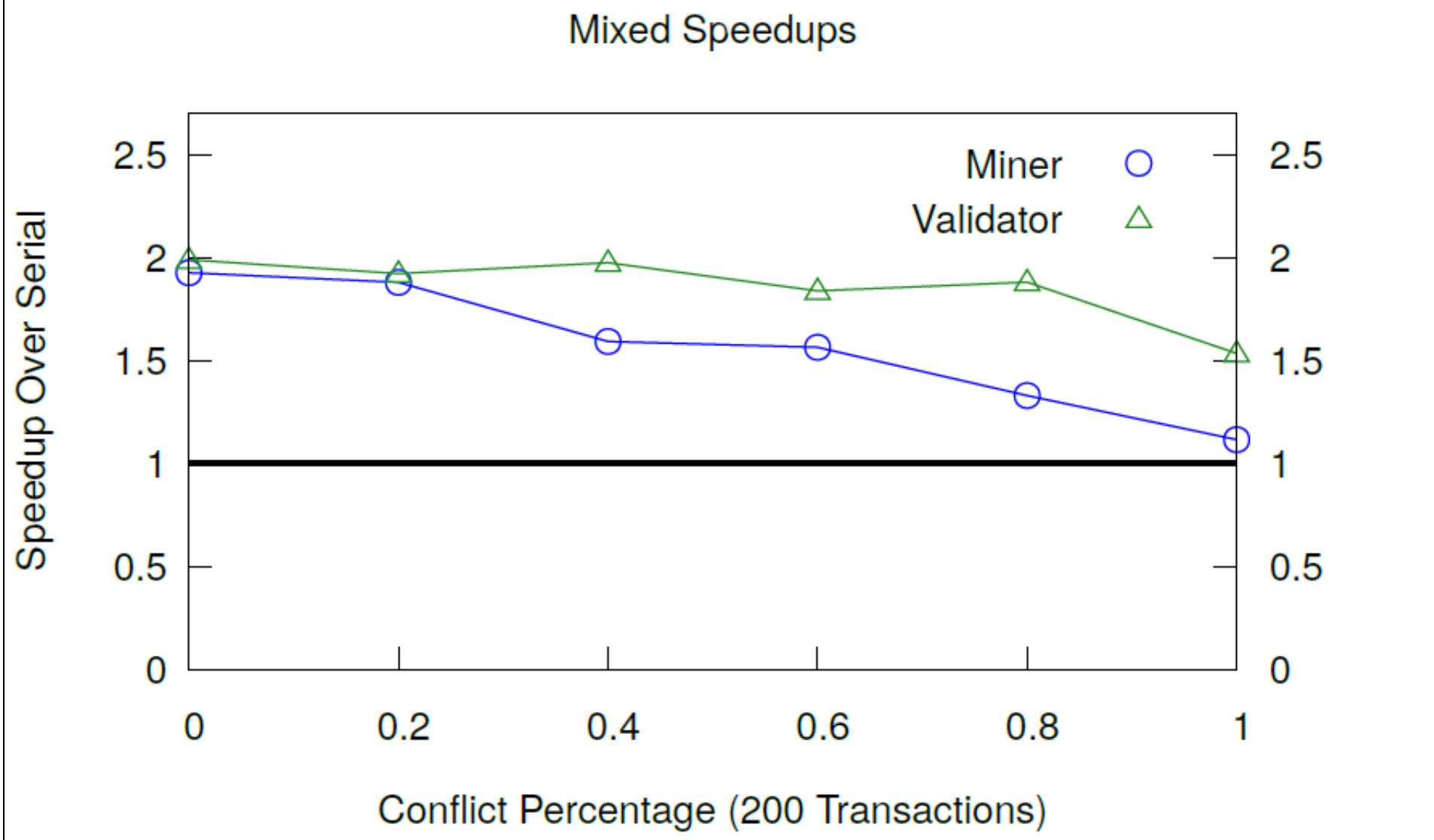
Benchmark #4: Mixed

All of the above

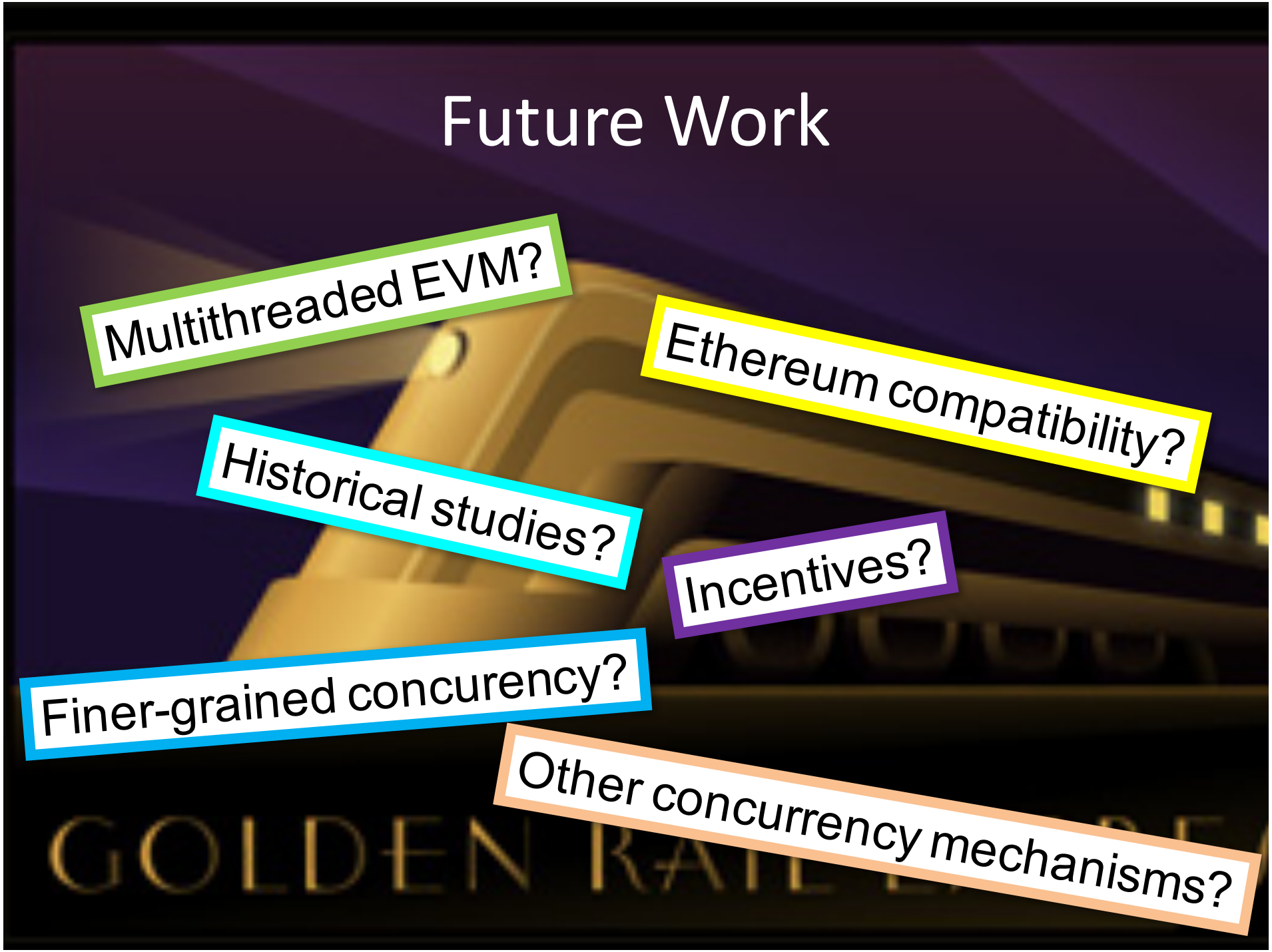
Equal proportions

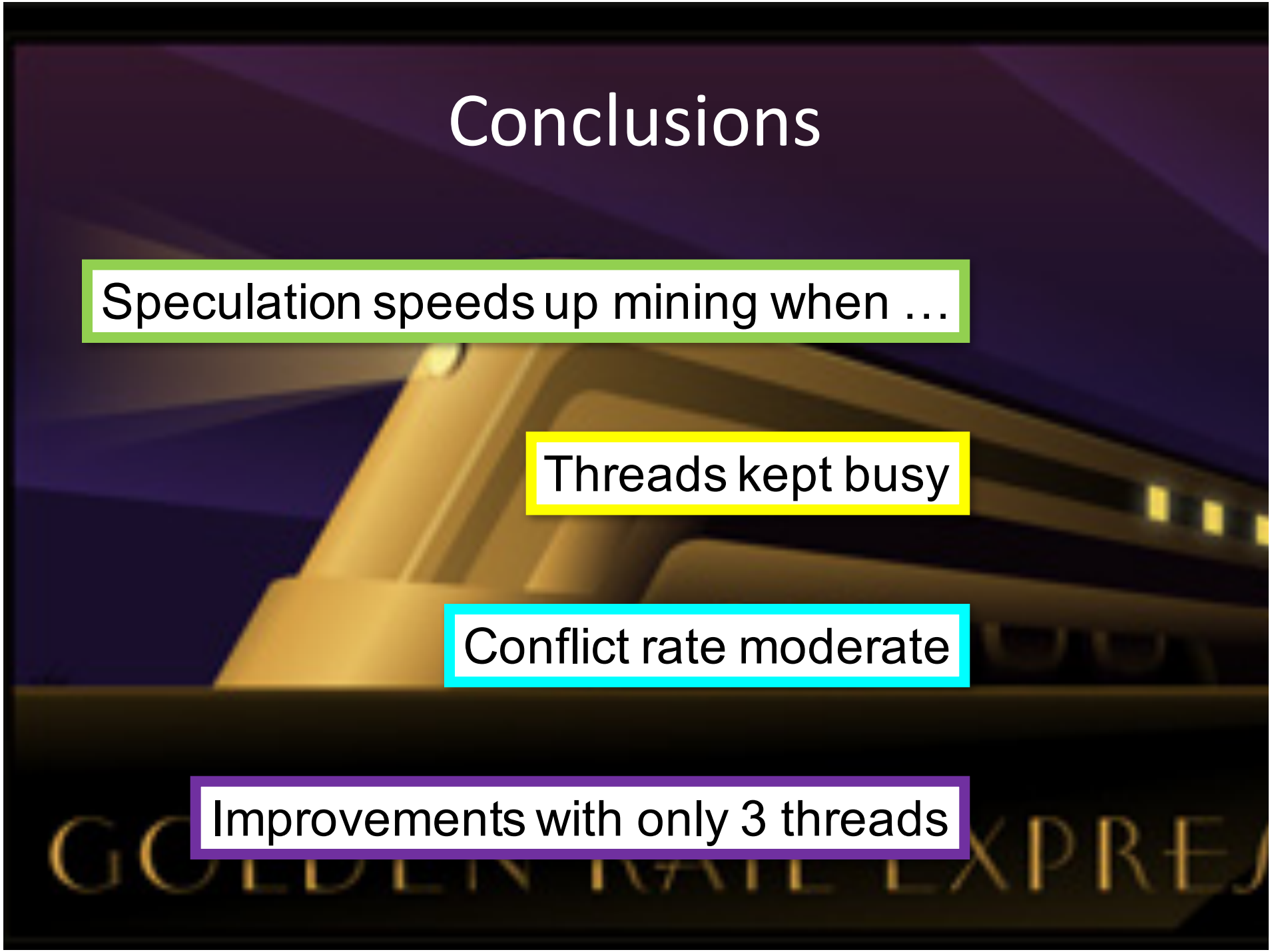


Varying Transactions per Block



Varying Levels of Conflict


Future Work
Multithreaded EVM?
Ethereum compatibility?
Historical studies?
Incentives?
Finer-grained concurency?
Other concurrency mechanisms?


Conclusions
Speculation speeds up mining when …
Threads kept busy
Conflict rate moderate
Improvements with only 3 threads


Thank You!
Questions?
GOLDEN RAIL EXPRE